

Introduction to Wavelet Trees

Giovanni Manzini

rank and select operations

Given a string $S[1,n]$ over an alphabet A we define:

$\text{rank}_S(c,i) = \#$ occurrences of symbol c in $S[1,i]$

$\text{select}_S(c,i) =$ position of the i -th c in $S[1,n]$

Example:

$S =$ abracadabra

$\text{rank}_S(a,3)=1, \quad \text{select}_S(a,3)=6$

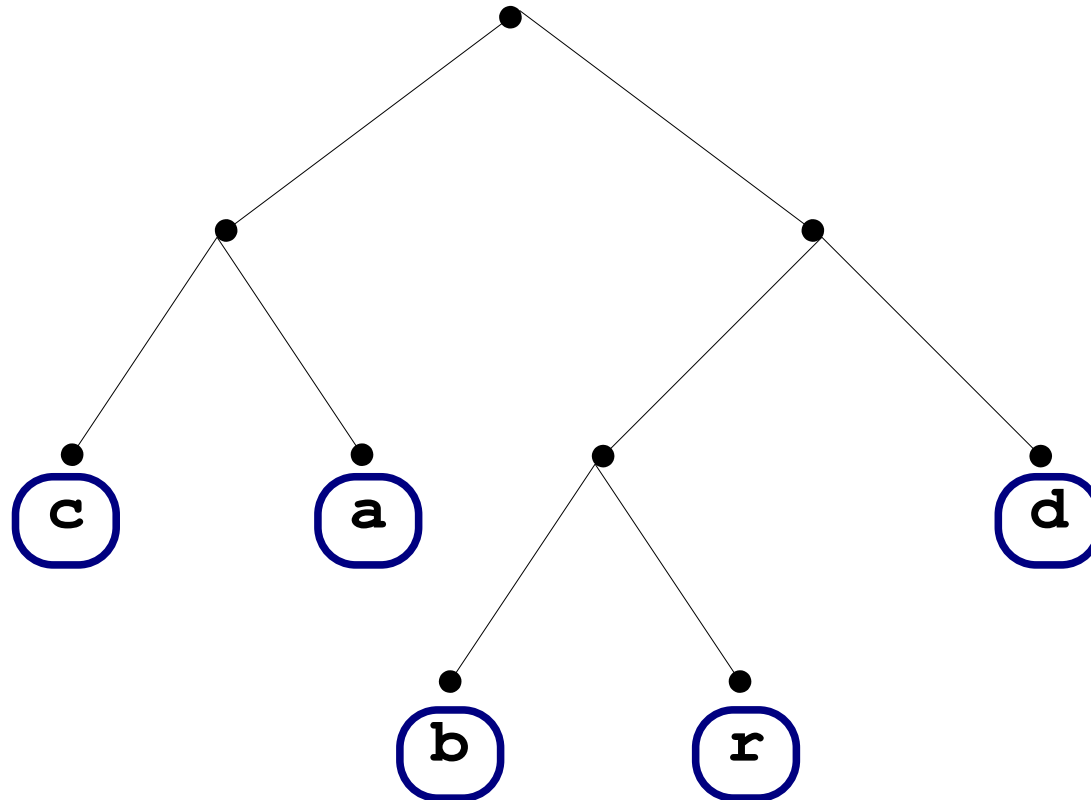
The Wavelet Tree data structure

Wavelet Trees have been introduced to represent compactly a string supporting rank/select operations.

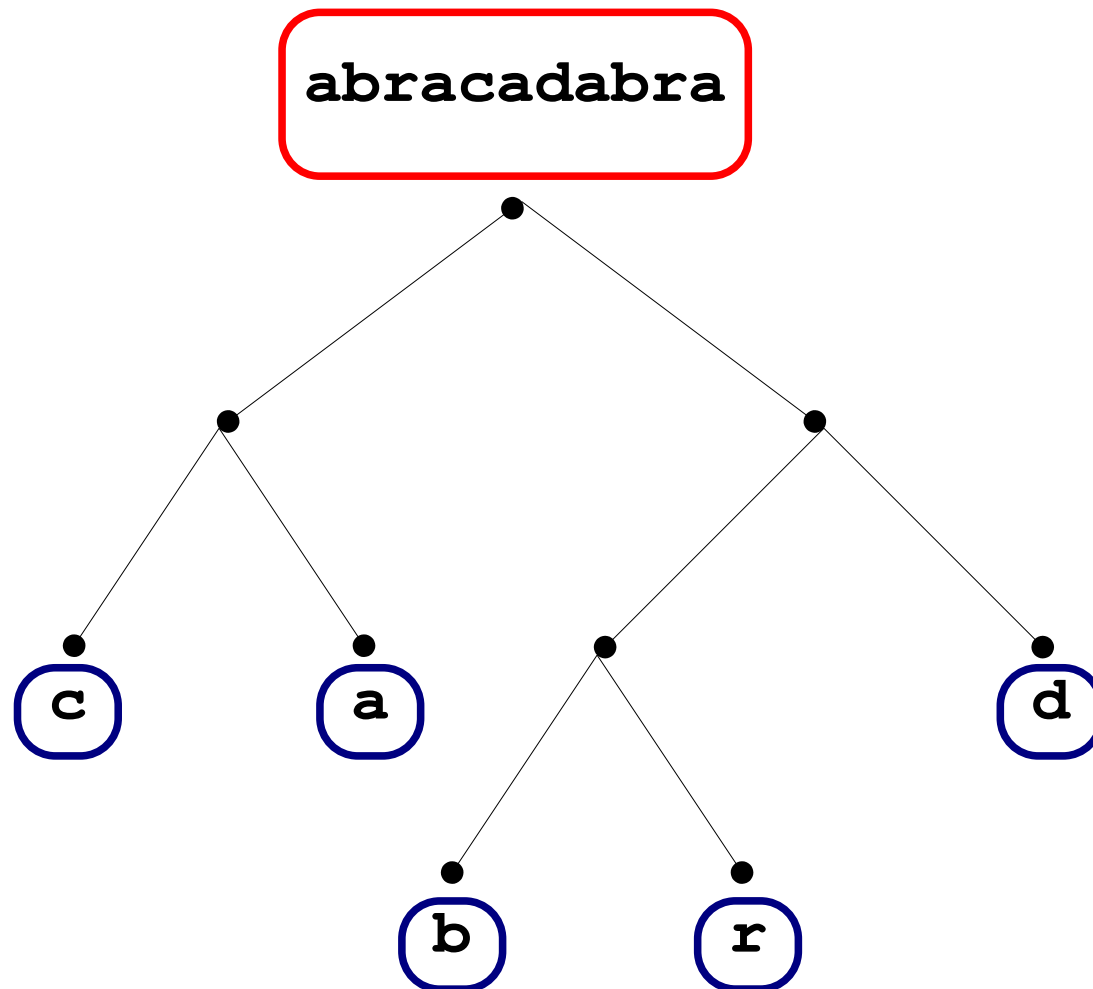
They work assuming we can do rank/select on binary strings,

To build a Wavelet Tree we start with a complete binary tree with a leaf for each alphabet symbol.

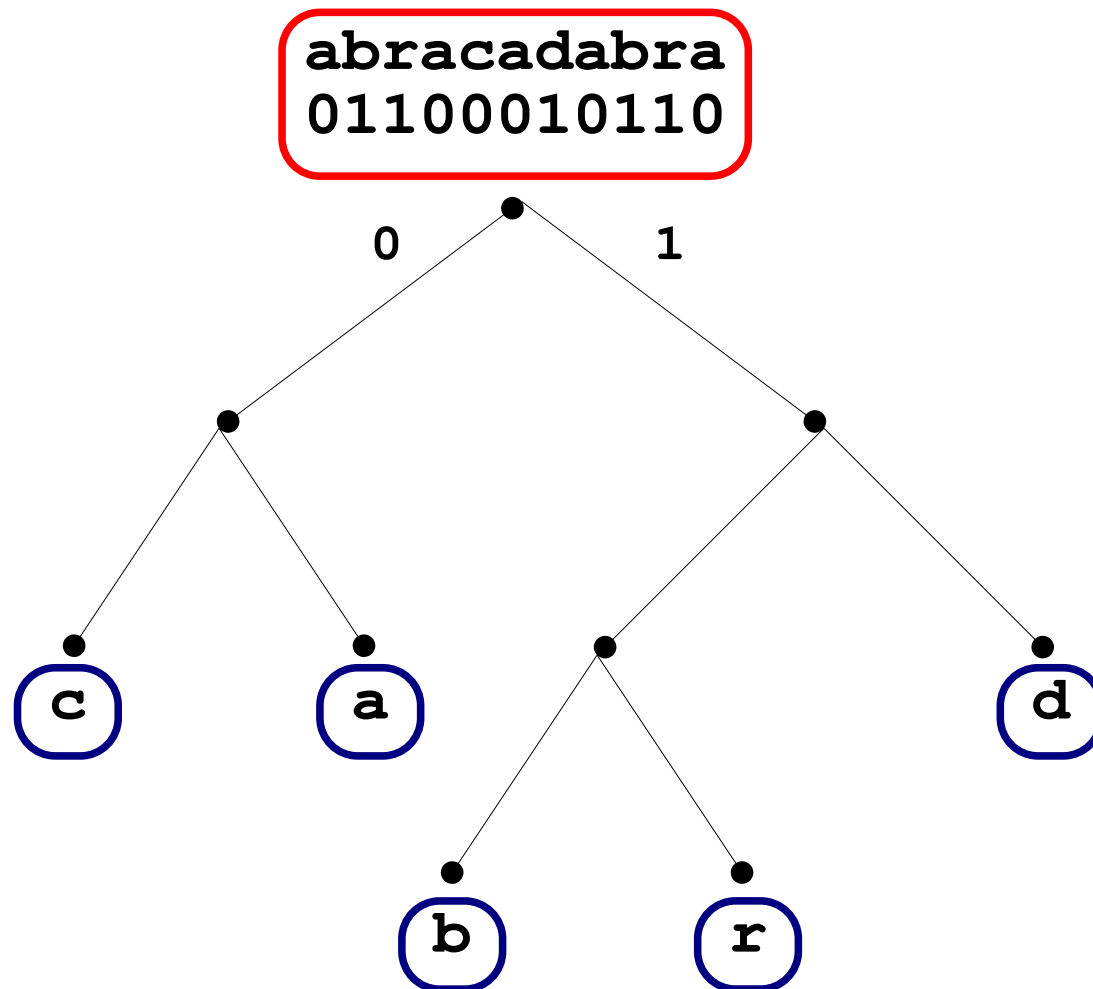
Example. $A = \{a, b, c, d, r\}$



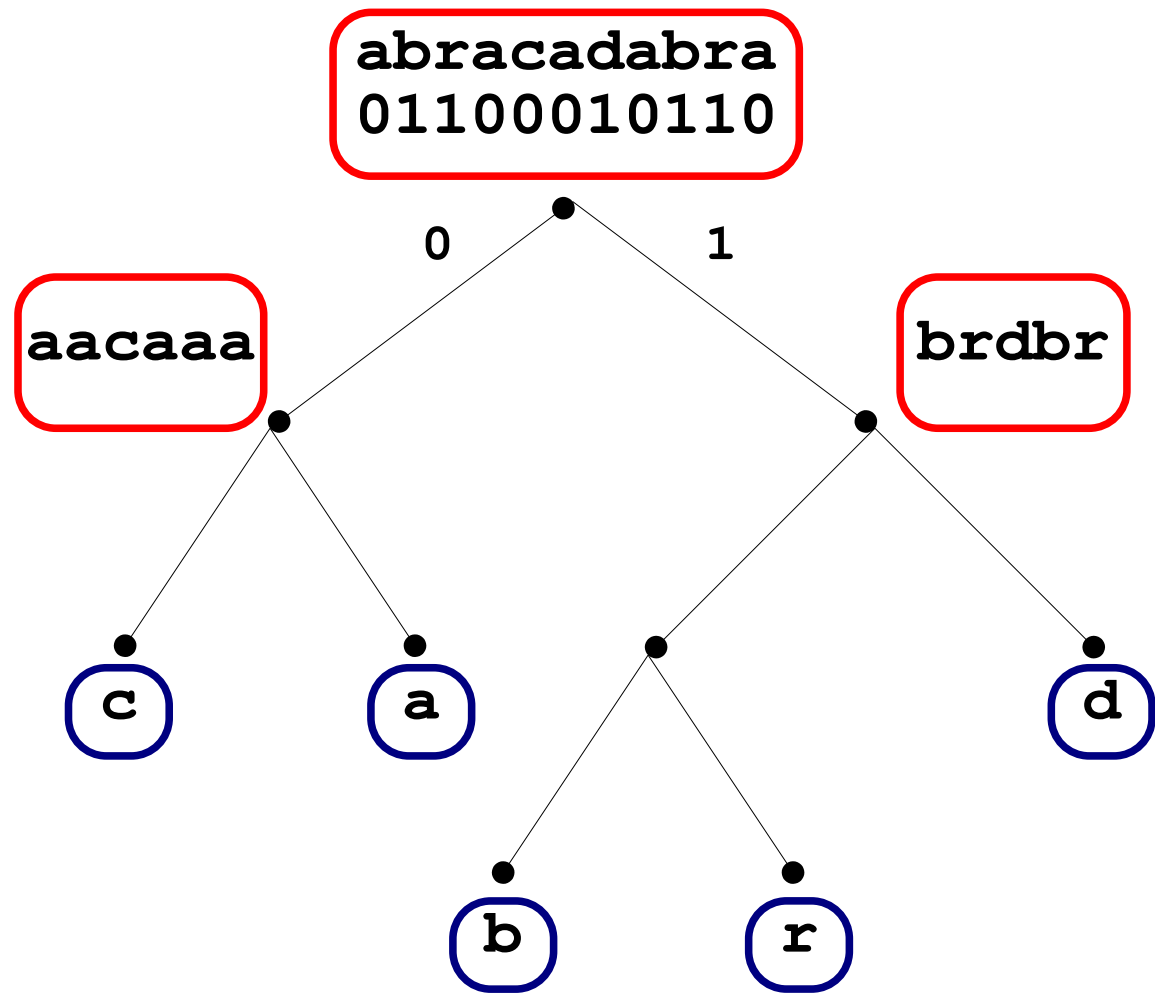
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



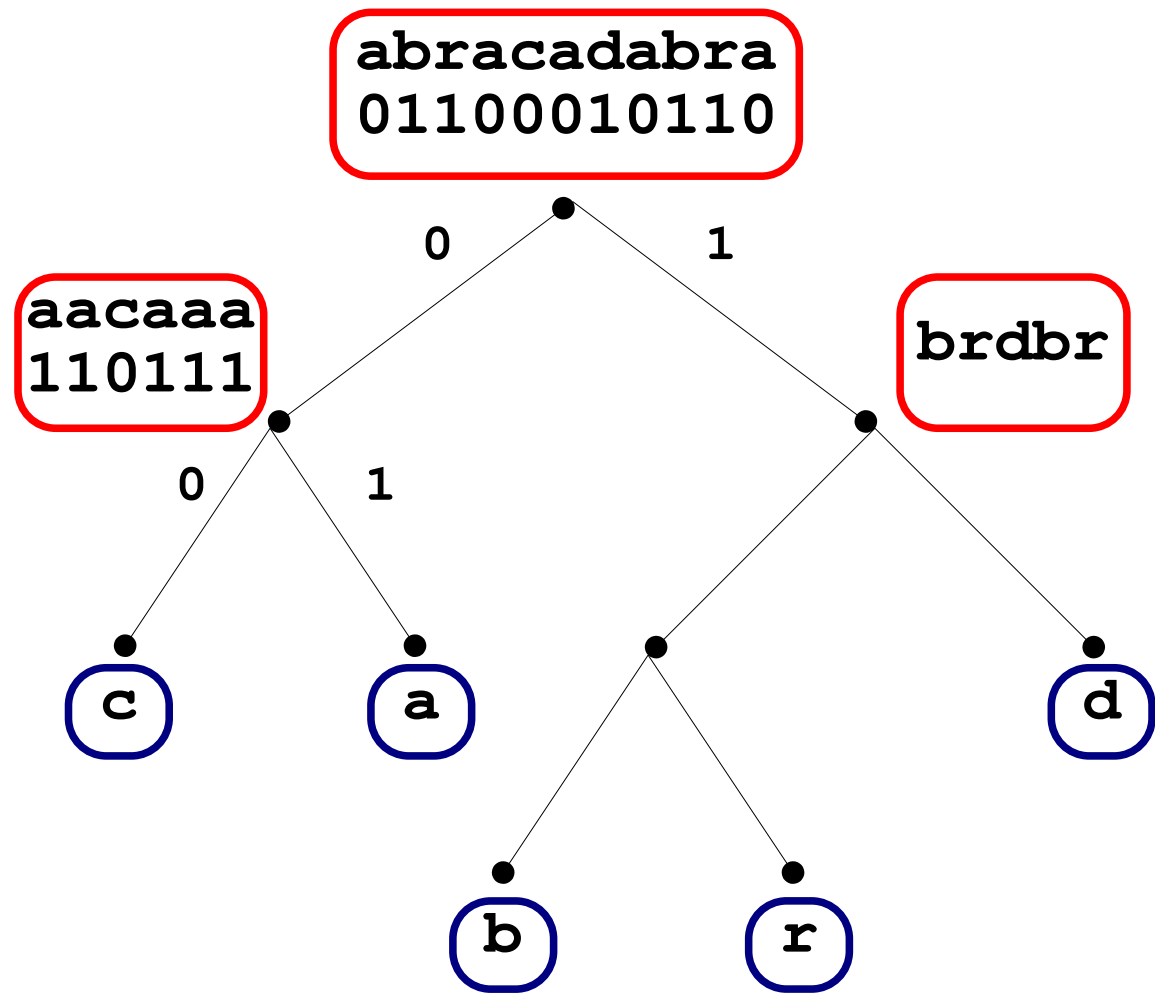
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



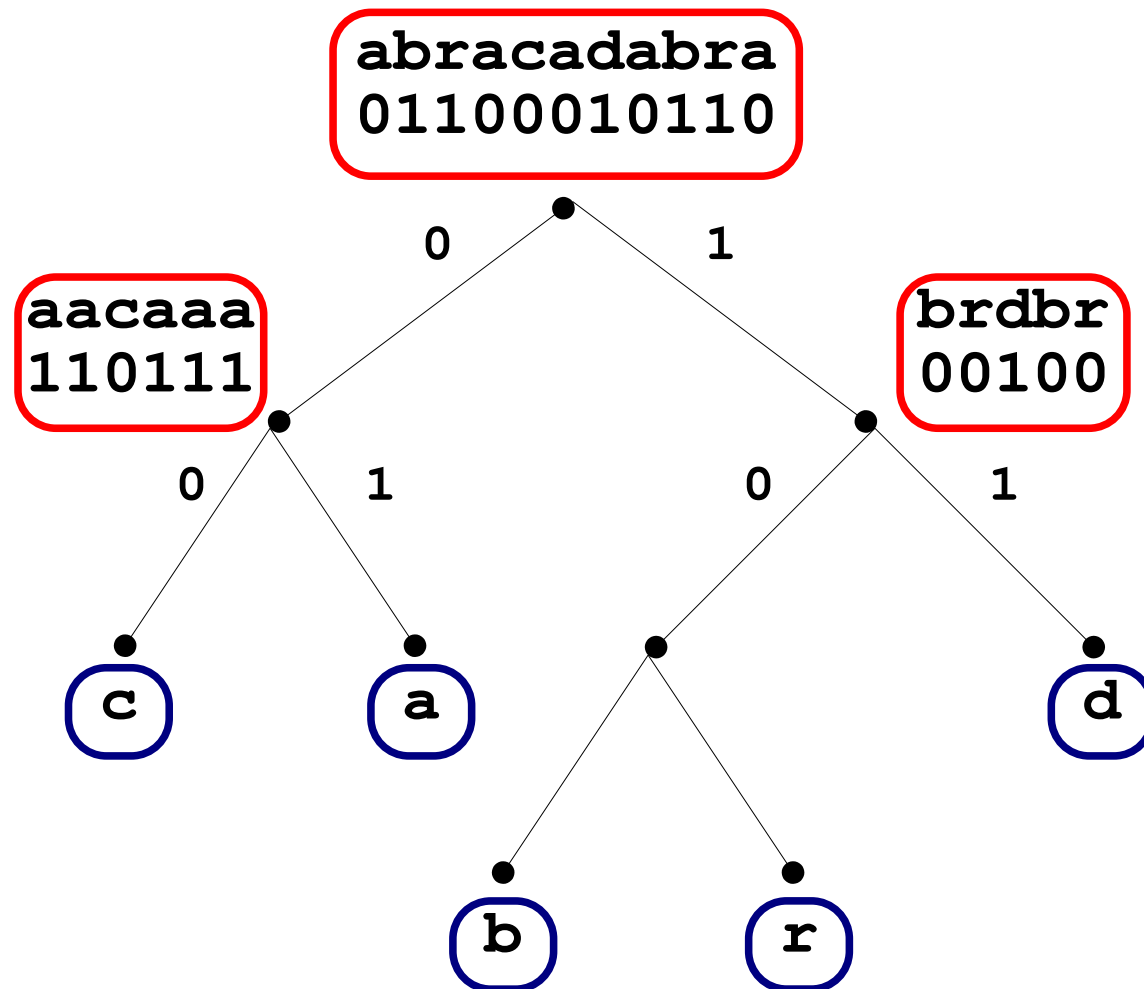
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



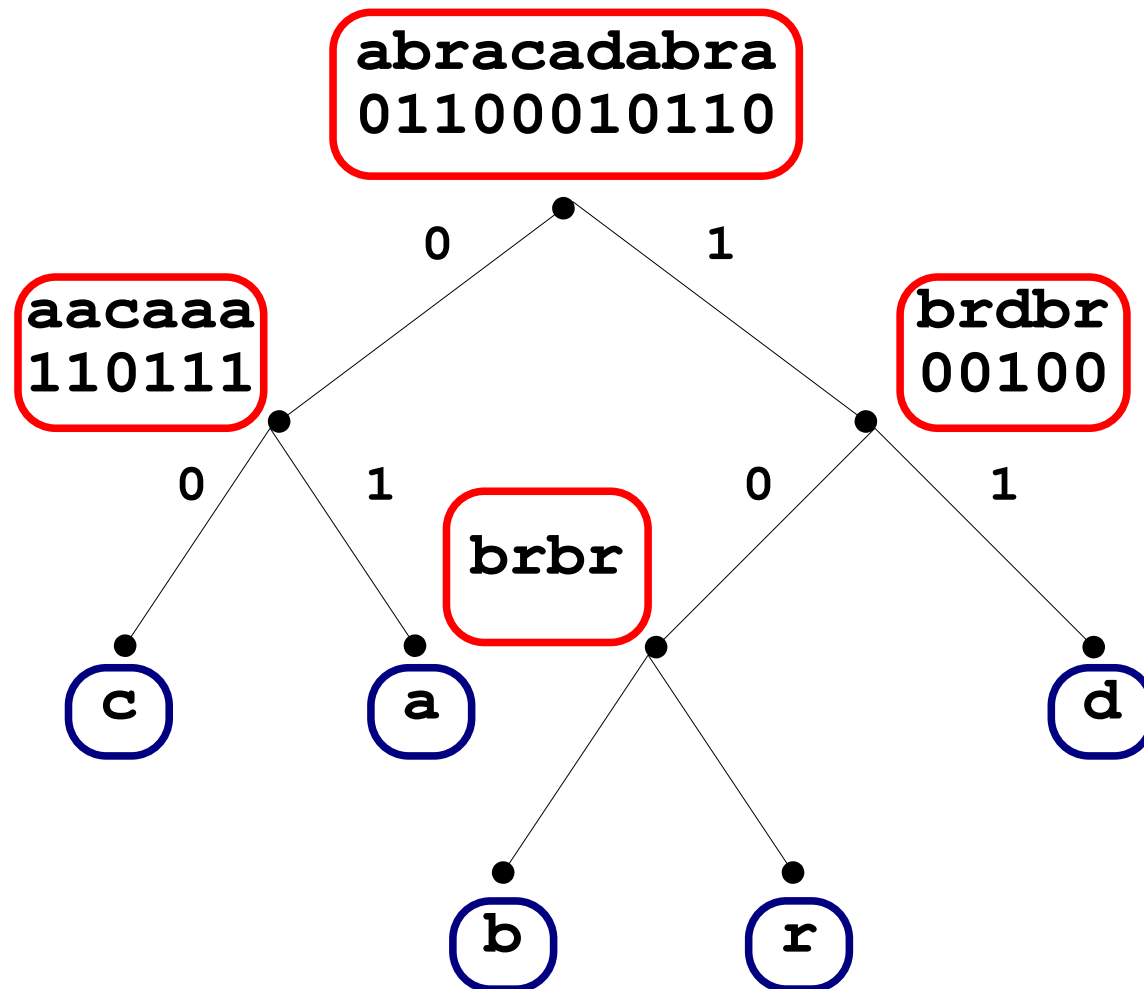
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



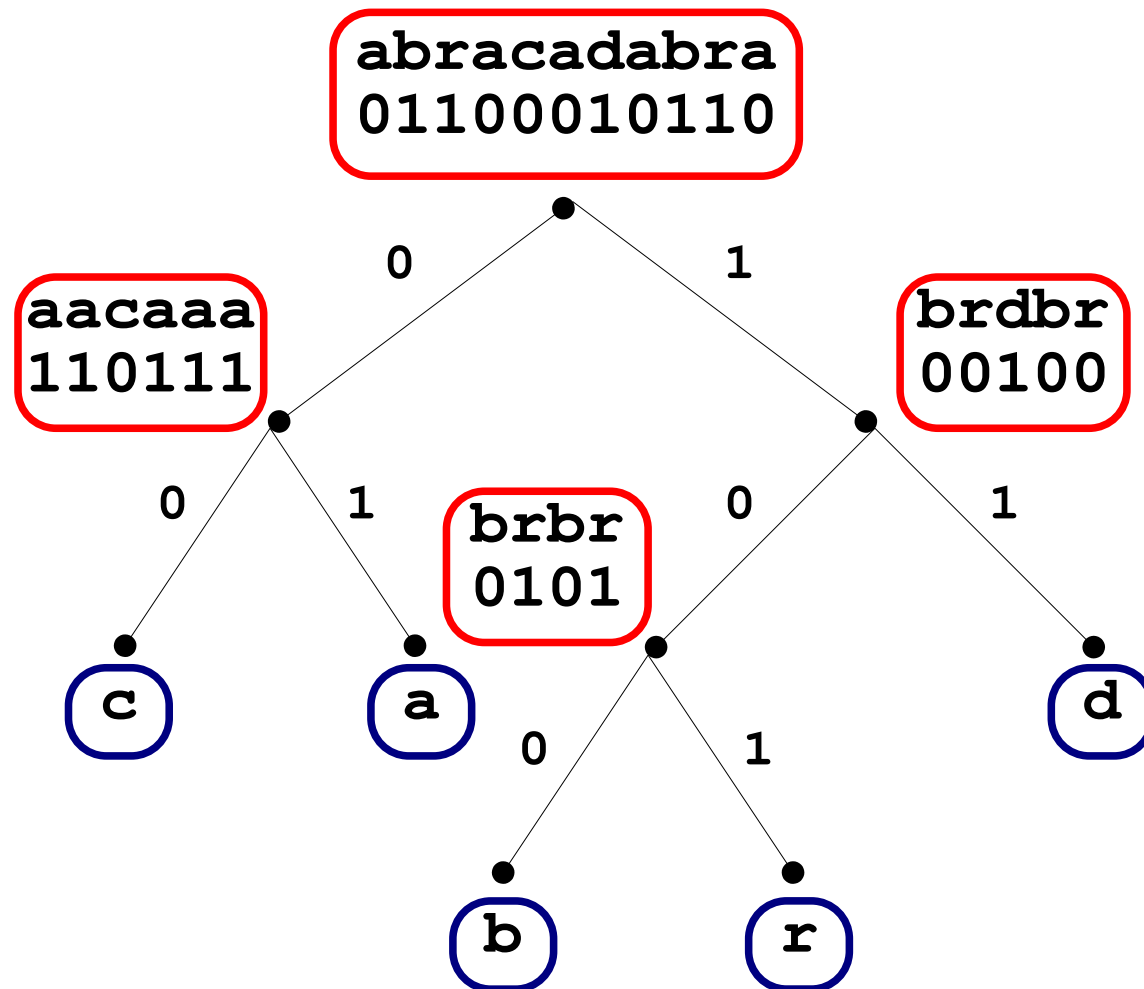
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



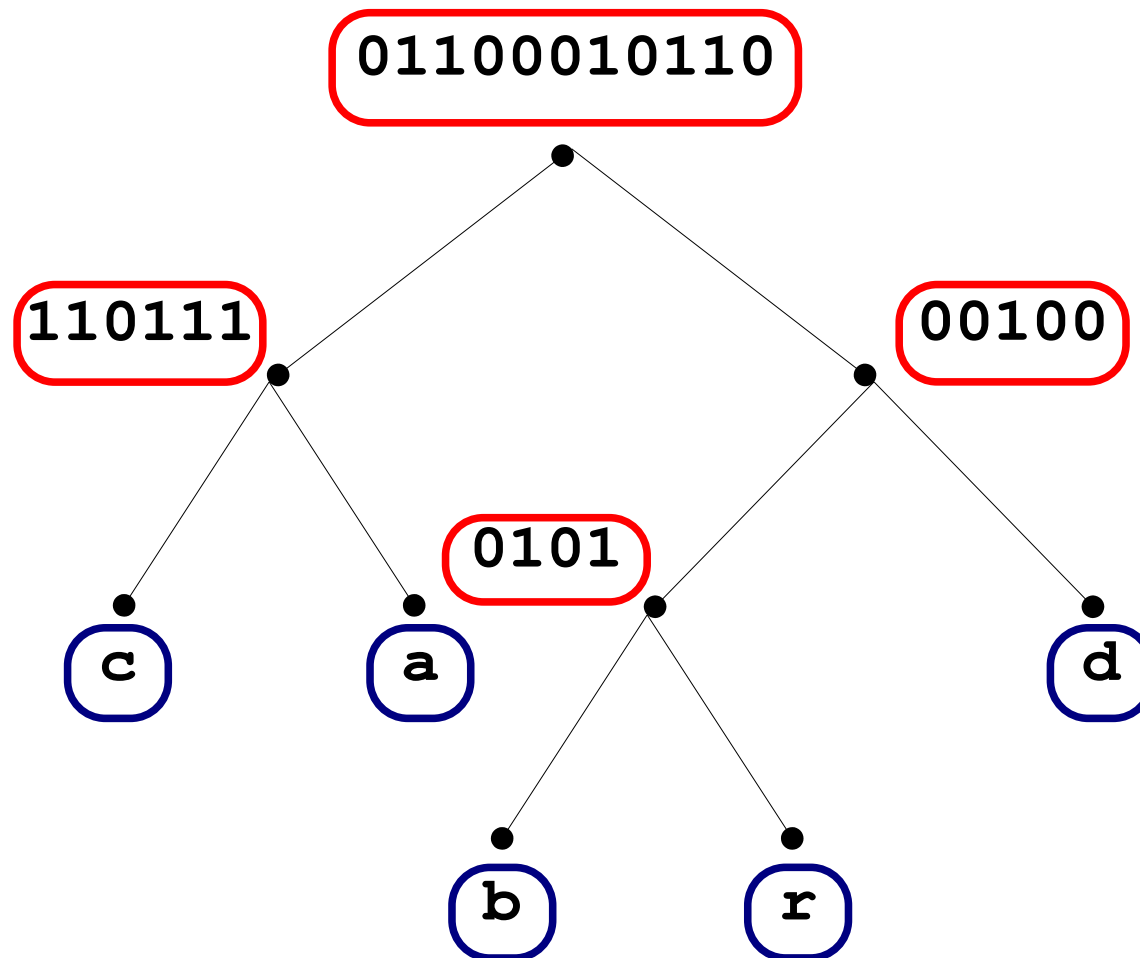
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



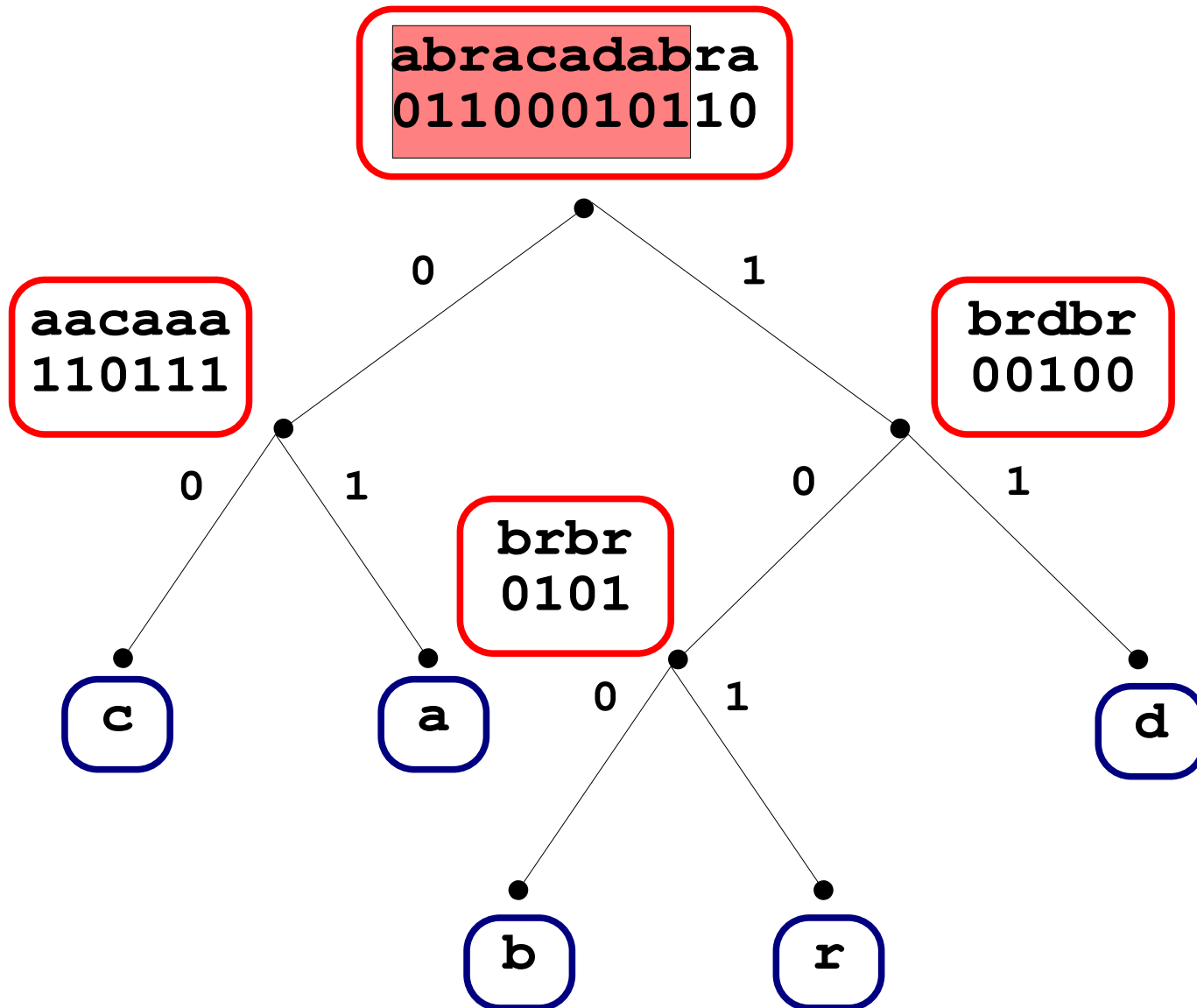
Given the string **abracadabra** we build the corresponding **Wavelet Tree** as follows:



rank/select queries over the original string can be answered via rank/select queries over the binary strings associated to internal nodes.



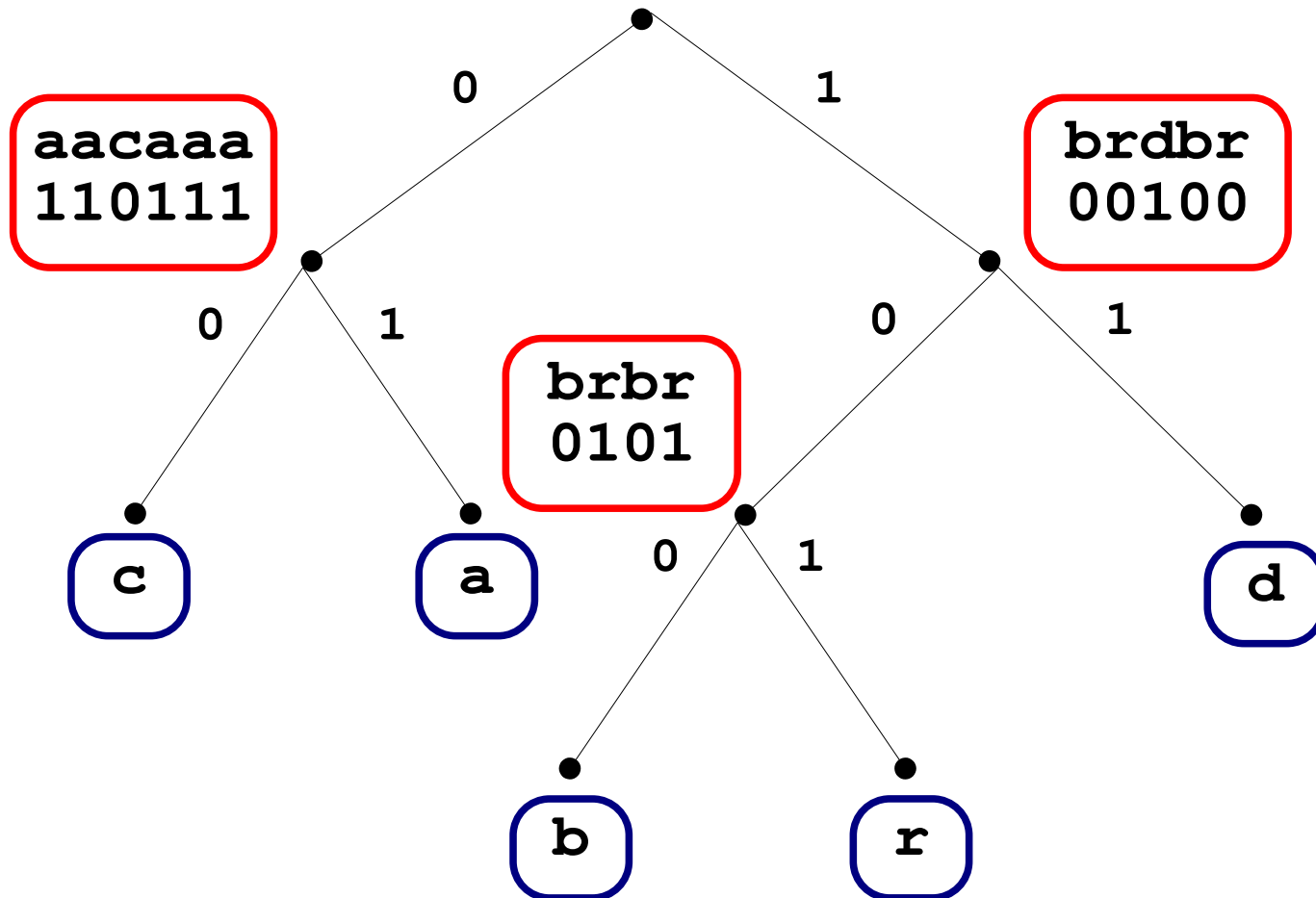
Suppose we want to compute $\text{rank}(b, 9)$



Suppose we want to compute $\text{rank}(b,9)$

abracadabra
01100010110

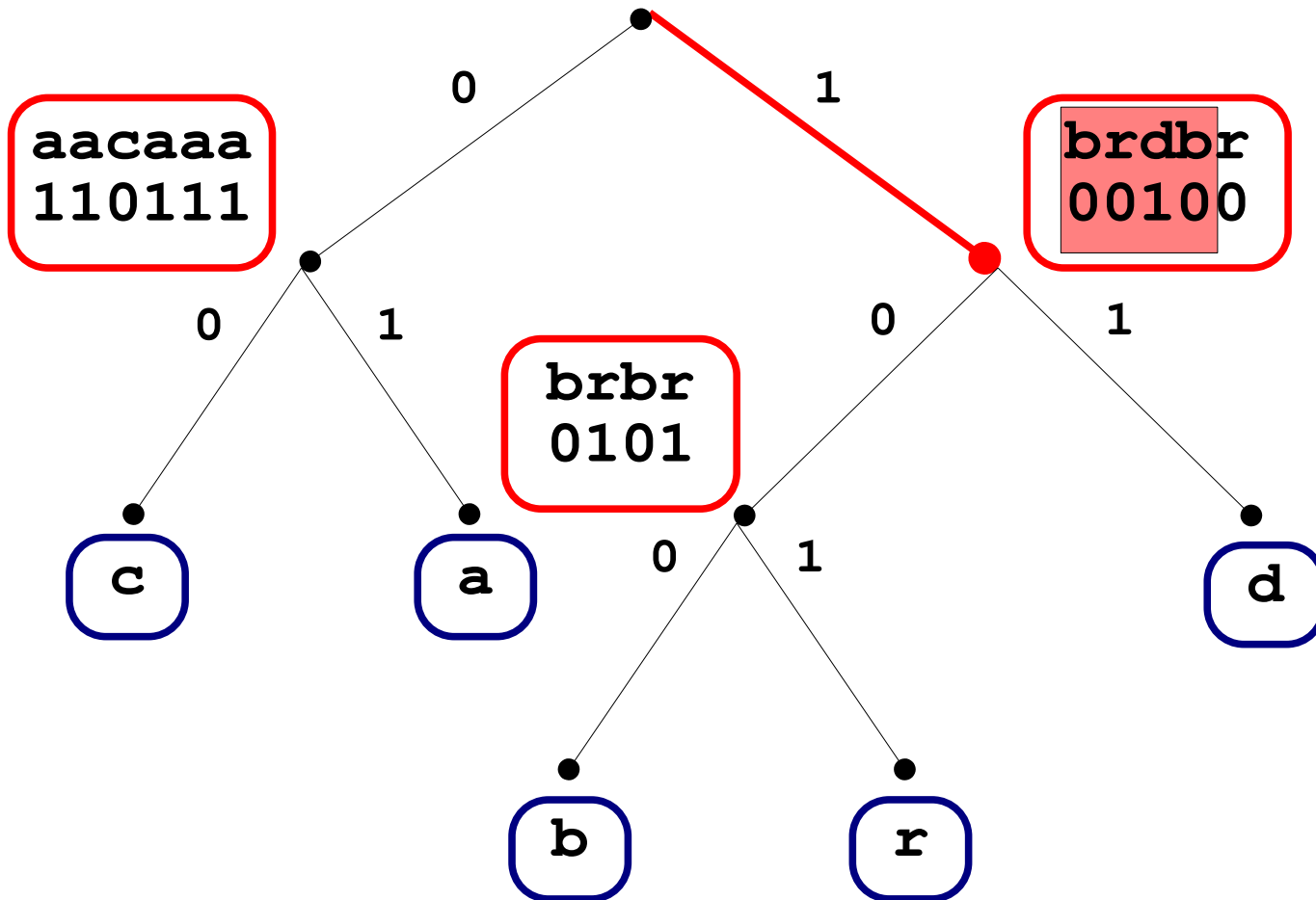
$$\text{rk}_1(9) = \text{rk}(b,9) + \text{rk}(r,9) + \text{rk}(d,9) = 4$$



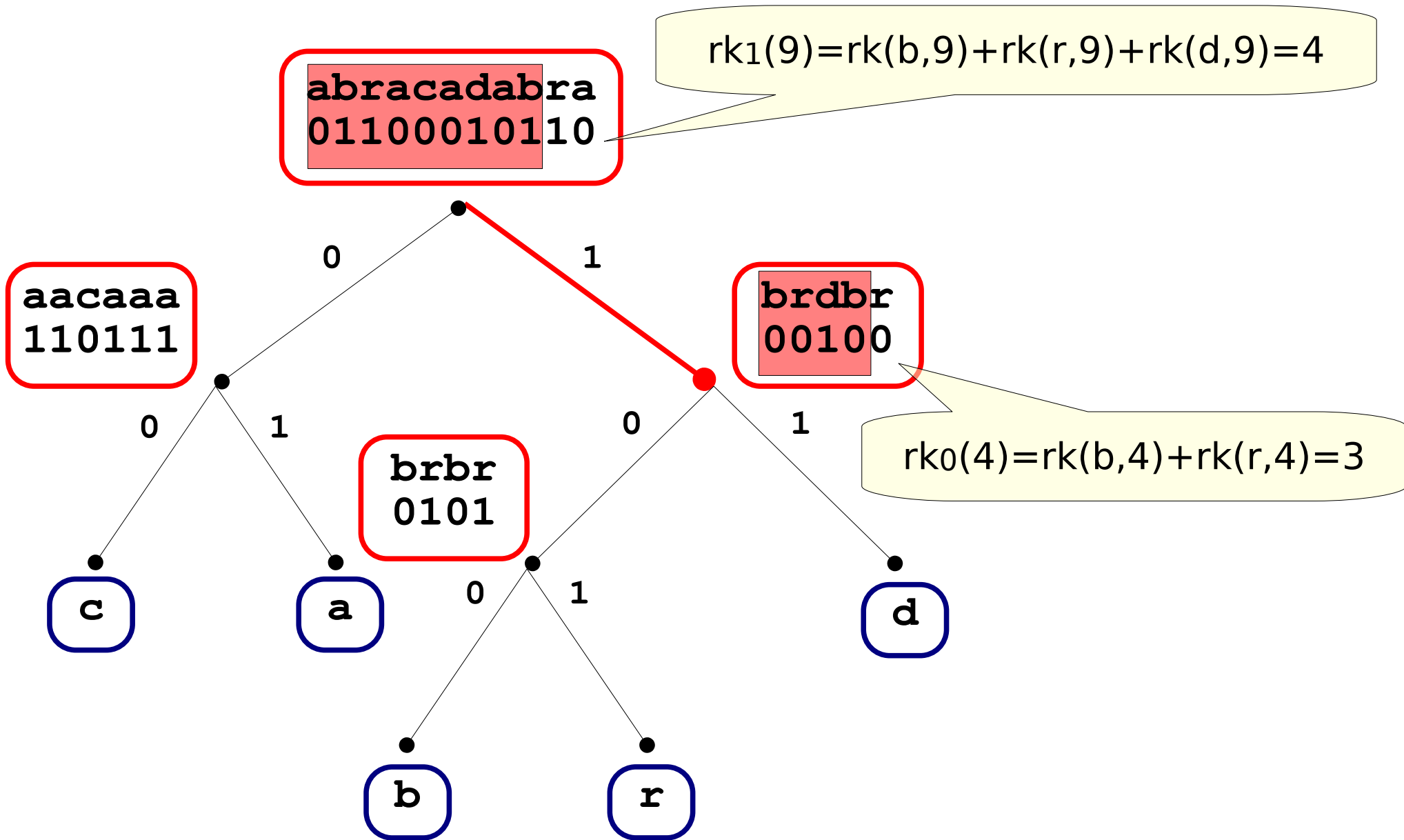
Suppose we want to compute $\text{rank}(b,9)$

abracadabra
01100010110

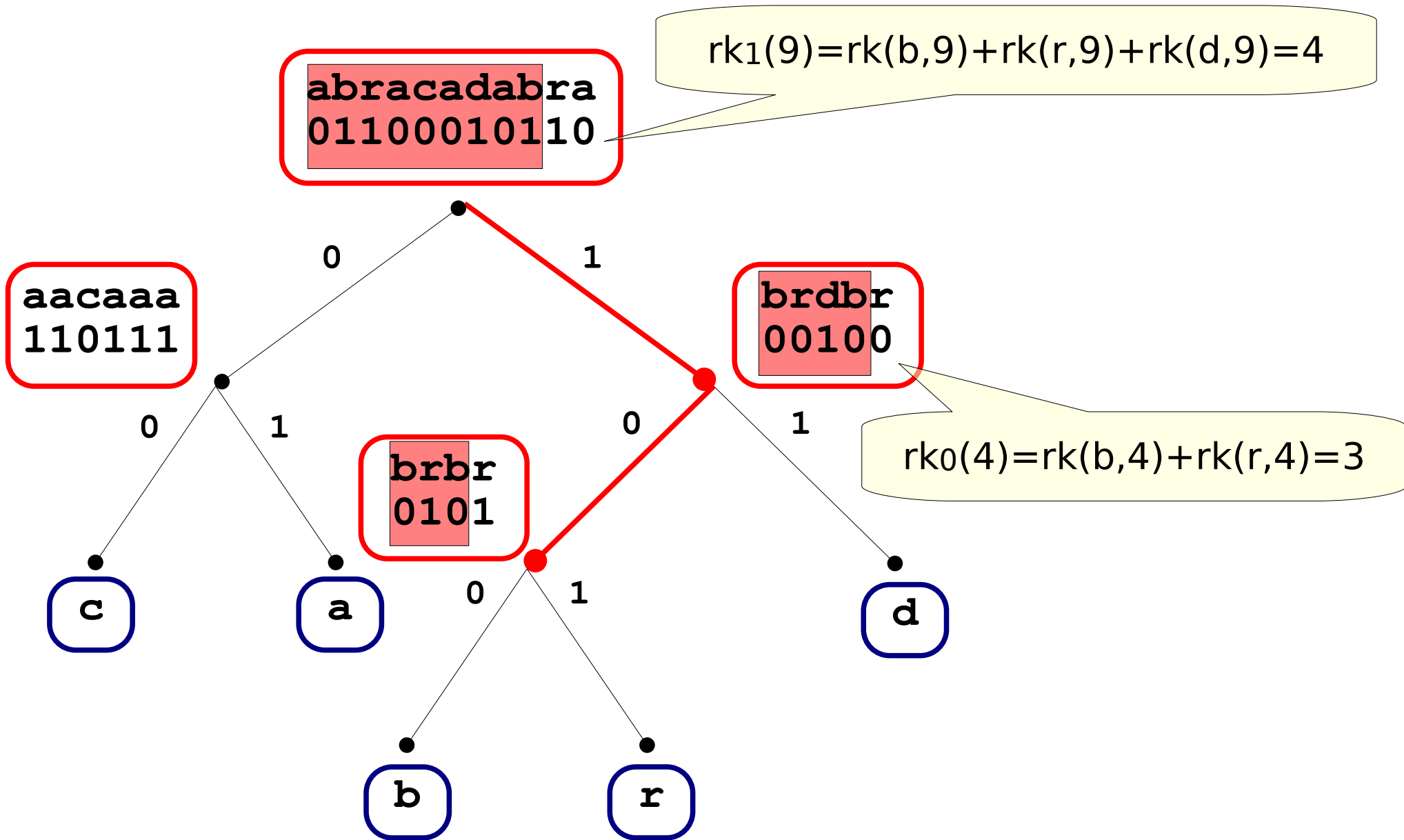
$$\text{rk}_1(9) = \text{rk}(b,9) + \text{rk}(r,9) + \text{rk}(d,9) = 4$$



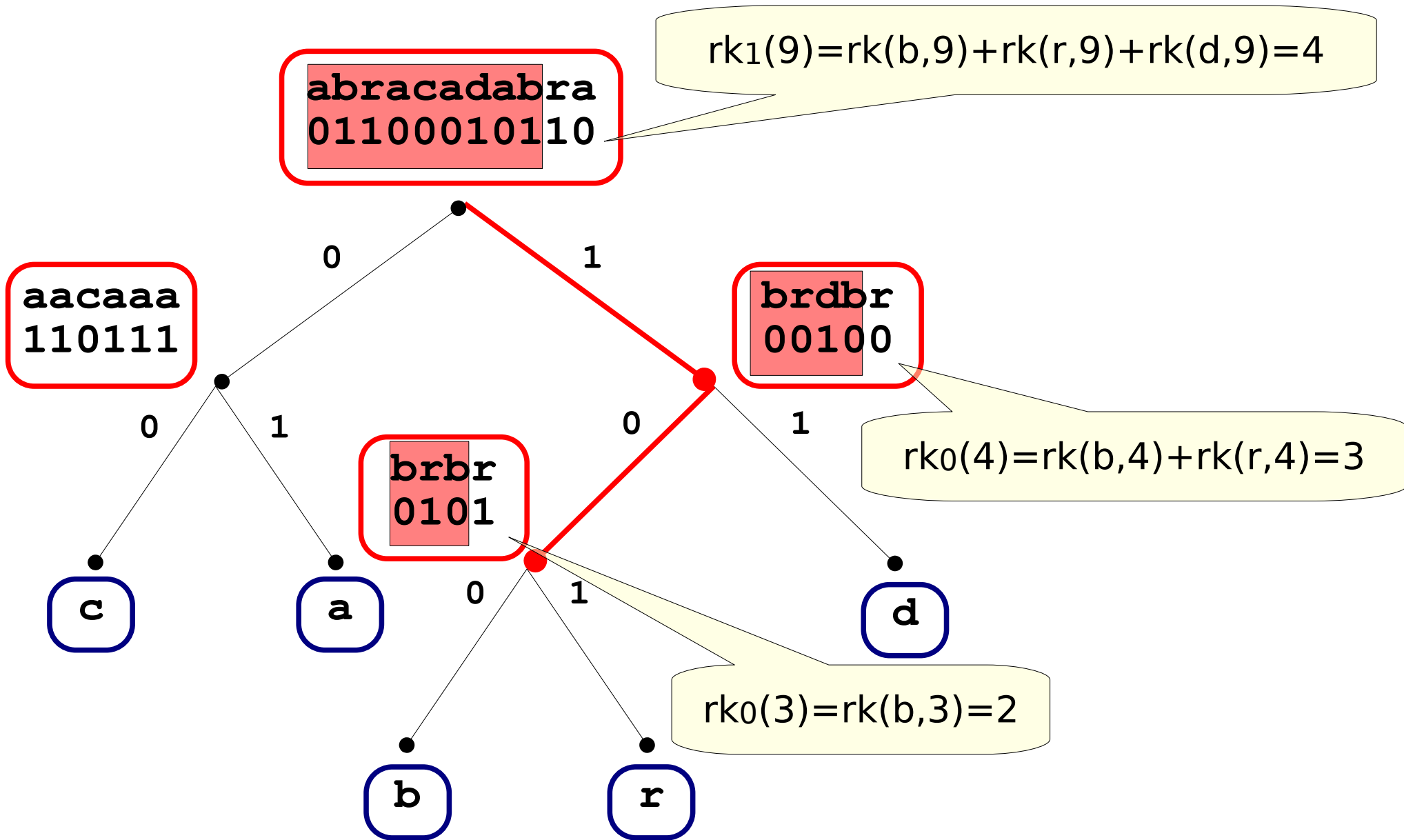
Suppose we want to compute $\text{rank}(b,9)$



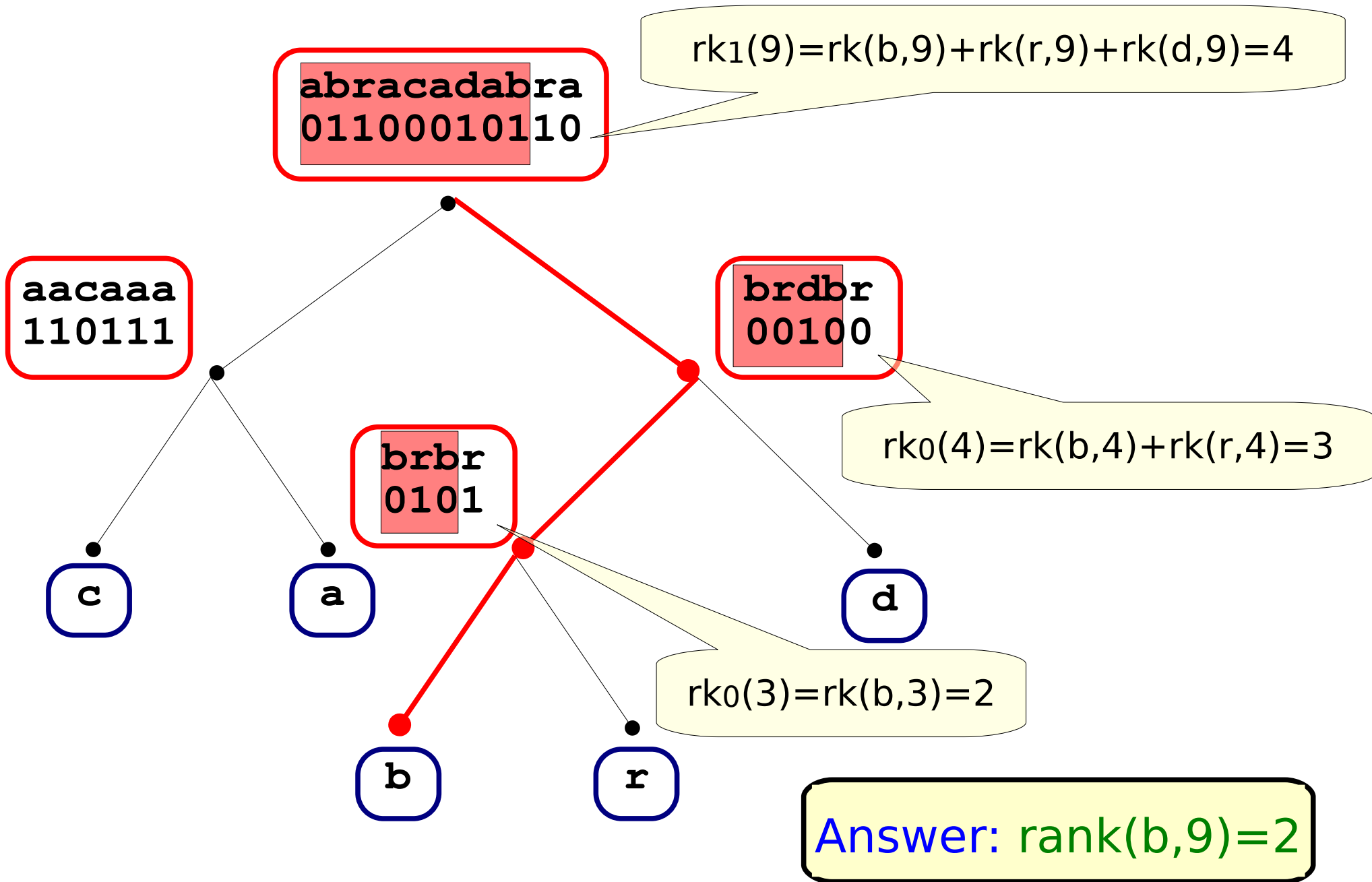
Suppose we want to compute $\text{rank}(b,9)$



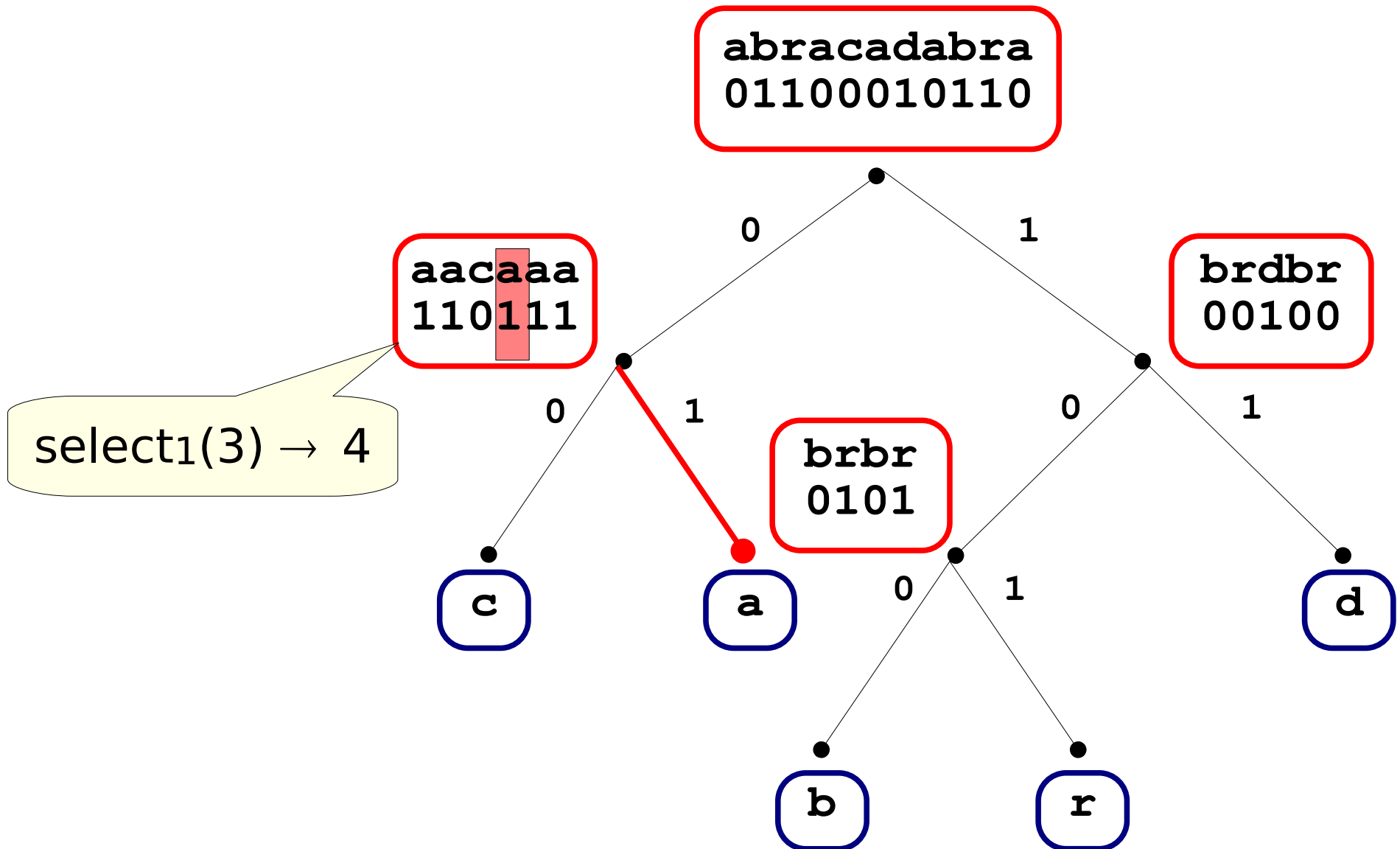
Suppose we want to compute $\text{rank}(b,9)$



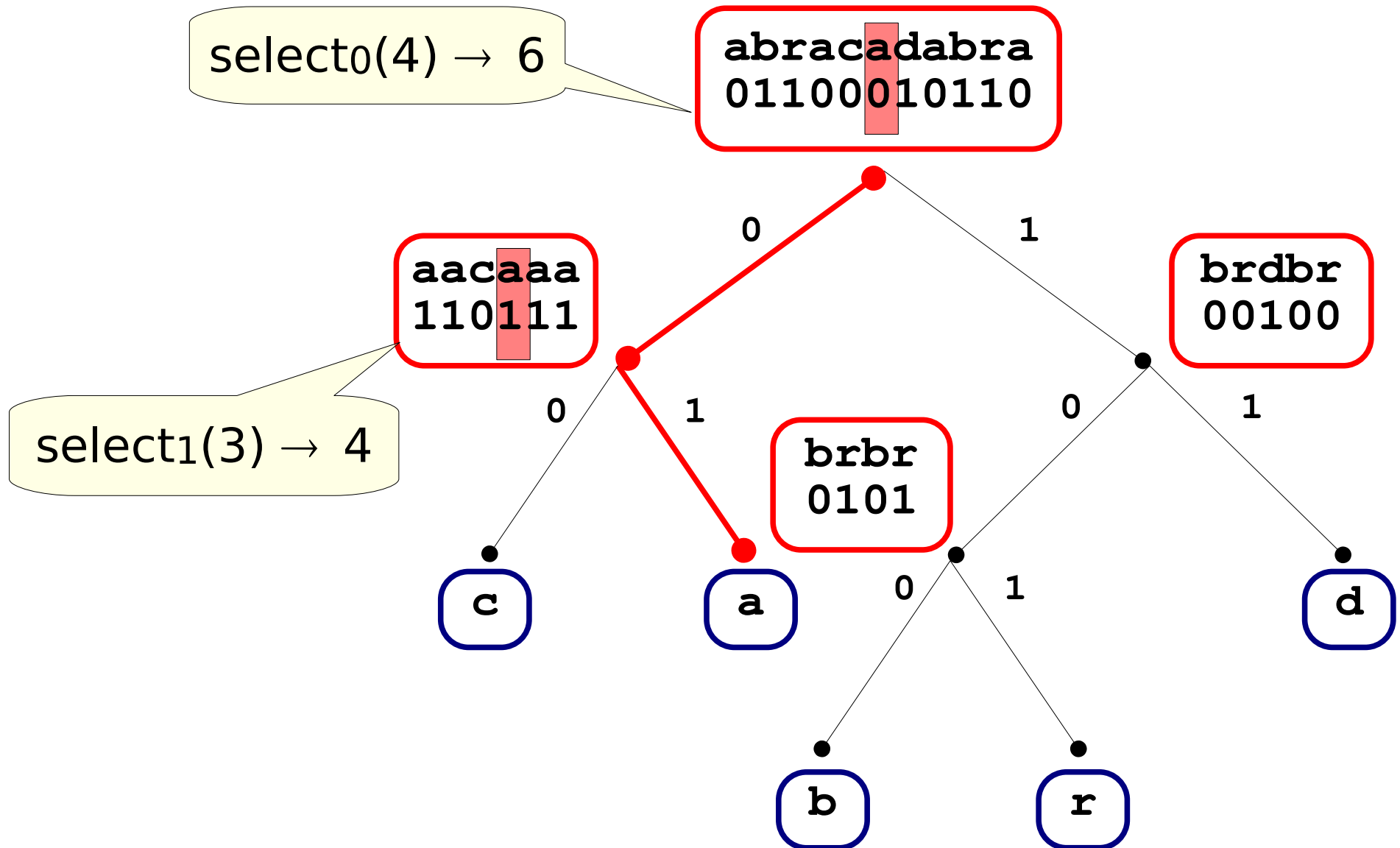
Suppose we want to compute $\text{rank}(b,9)$



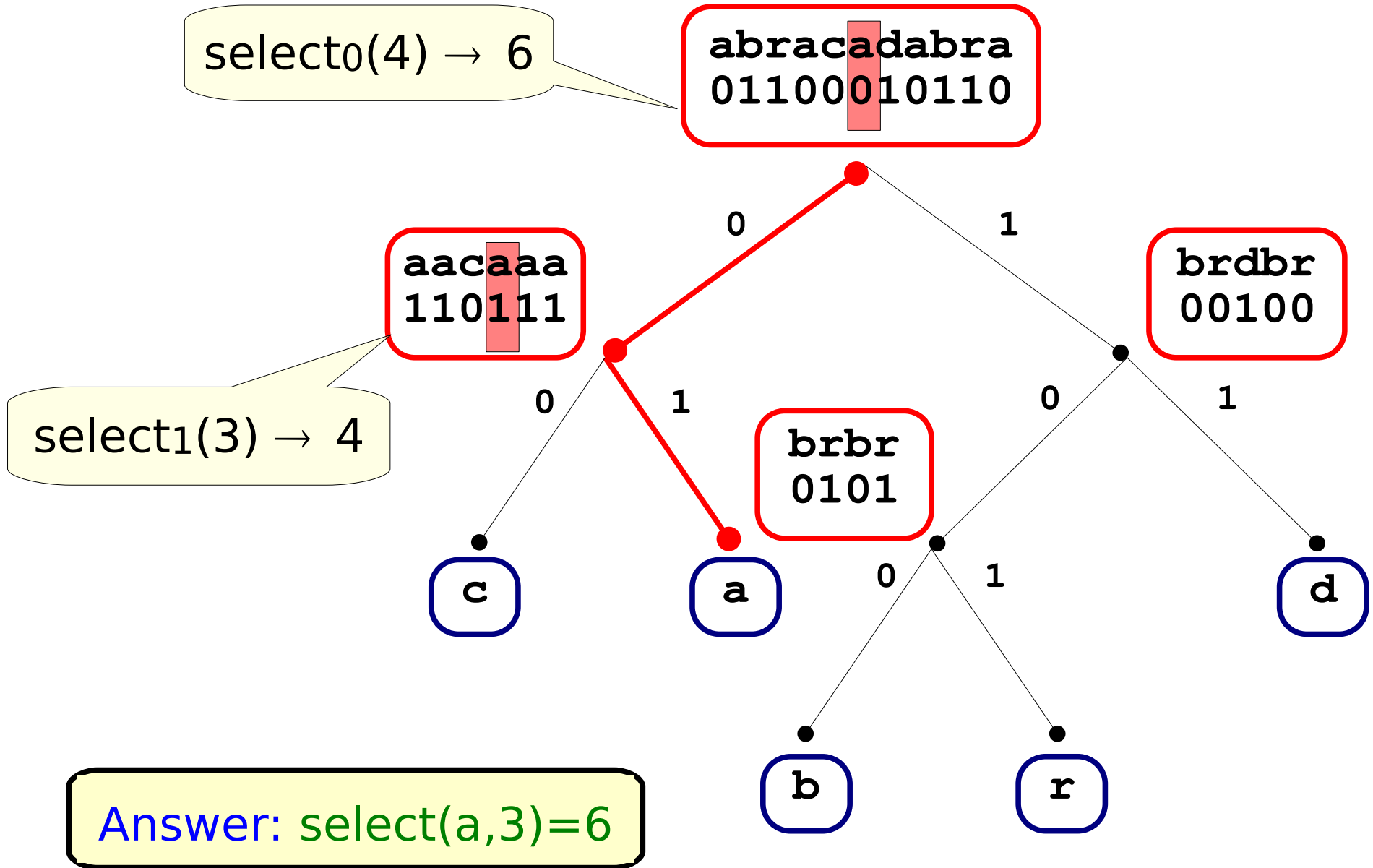
Select queries go upward. To compute $\text{select}(a,3)$



Select queries go upward. To compute $\text{select}(a,3)$

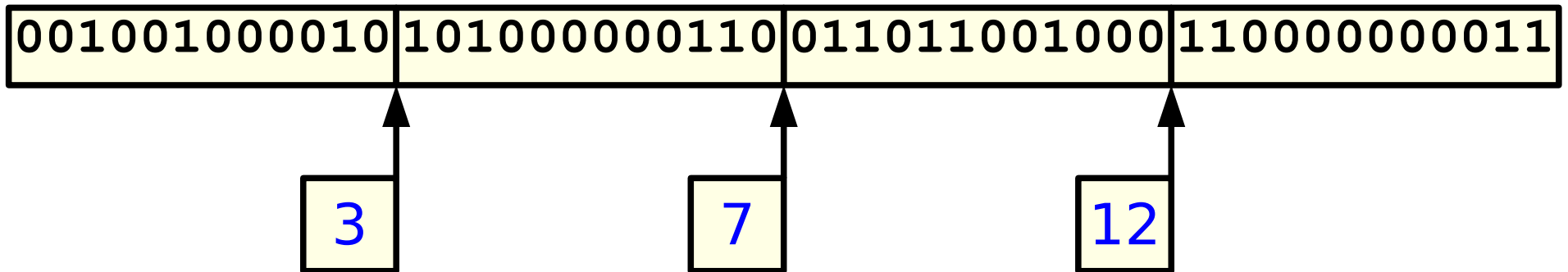


Select queries go upward. To compute $\text{select}(a,3)$

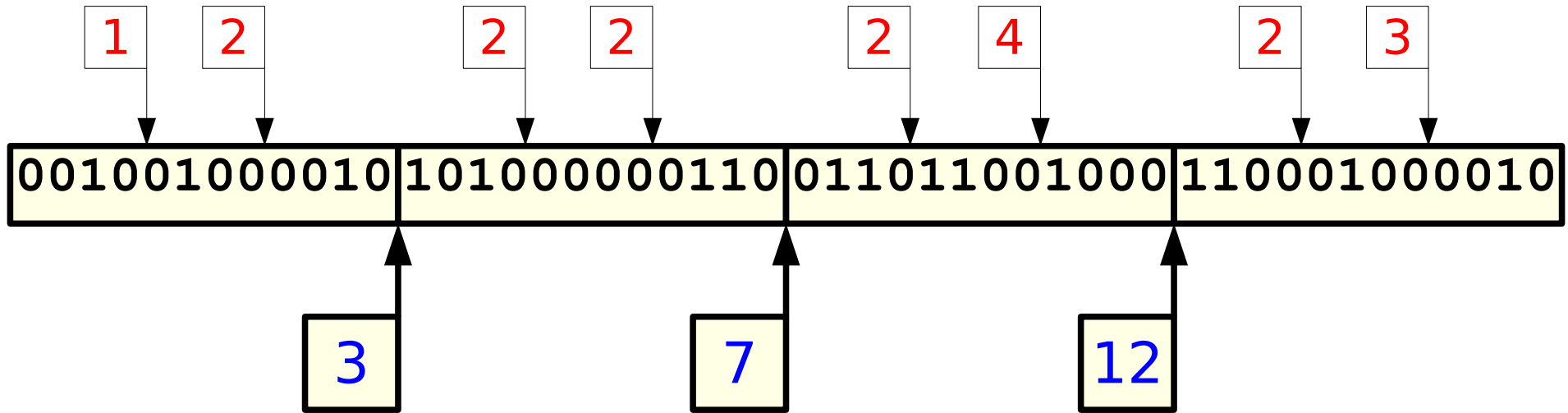


Rank₁ queries on binary strings

Simplest idea: split the binary string into blocks and store a partial sum for each block:

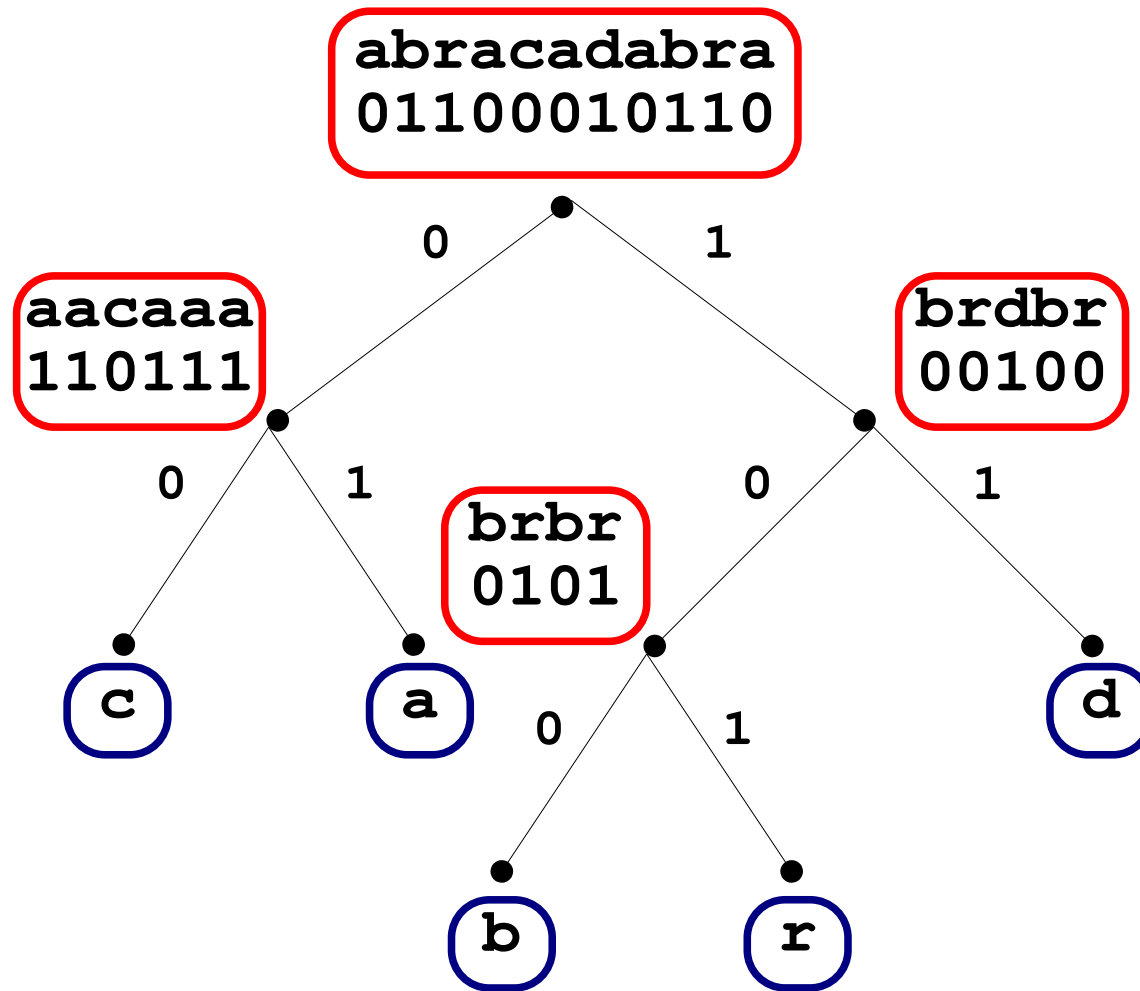


We can refine the idea using **blocks** and **mini-blocks**:



Similar techniques for select queries!

Wavelet Trees as scrambled prefix codes



We are implicitly using the encoding:

$c \rightarrow 00$ $a \rightarrow 01$ $b \rightarrow 100$ $r \rightarrow 101$ $d \rightarrow 11$

Wavelet Trees as scrambled prefix codes:

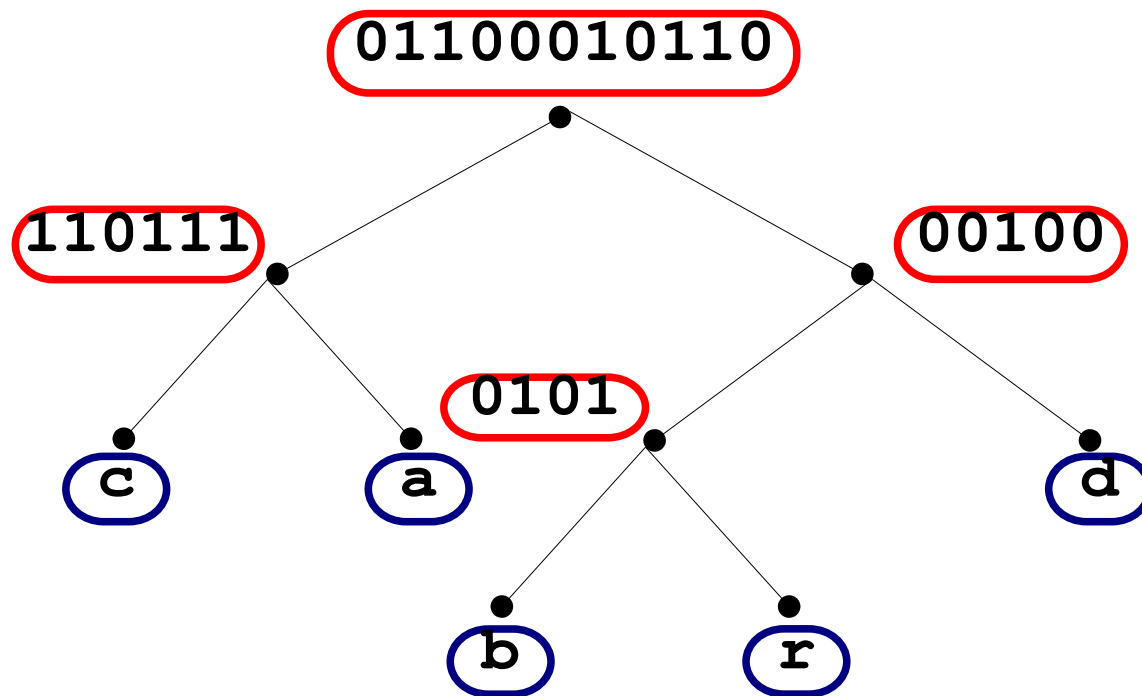
We are implicitly using the encoding:

$a \rightarrow 01$ $c \rightarrow 00$ $b \rightarrow 100$ $r \rightarrow 101$ $d \rightarrow 11$

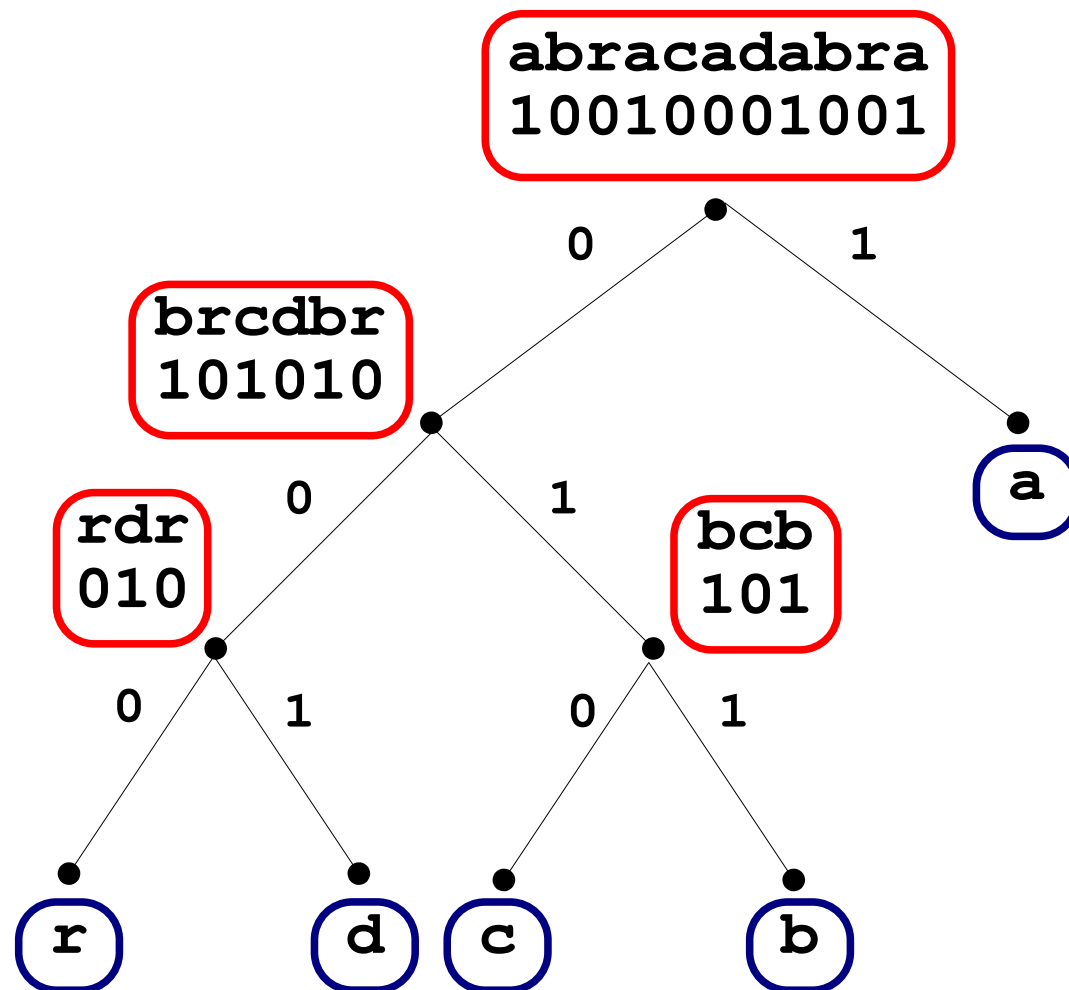
Traditional prefix codes:

abracadabra \rightarrow *01 100 101 01 00 01 11 ...*

Wavelets:



Changing the shape of the the Wavelet Tree,
all properties mentioned above still hold with
a different prefix code.



r → 000
d → 001
c → 010
b → 011
a → 1

We can choose **any** prefix-free binary encoding of the characters and build the corresponding **Wavelet Tree**.

If we choose **Huffman codes** we have compression for free and direct access to arbitrary positions in the text.

Summing up

Binary strings
operations

+

Wavelet
Trees

+

=

Compression and efficient rank support

+

BWT &
backward search

=

Compressed full text index

Other Wavelet Trees virtues

Simple solutions for range queries on an integer sequence $S[1,n]$

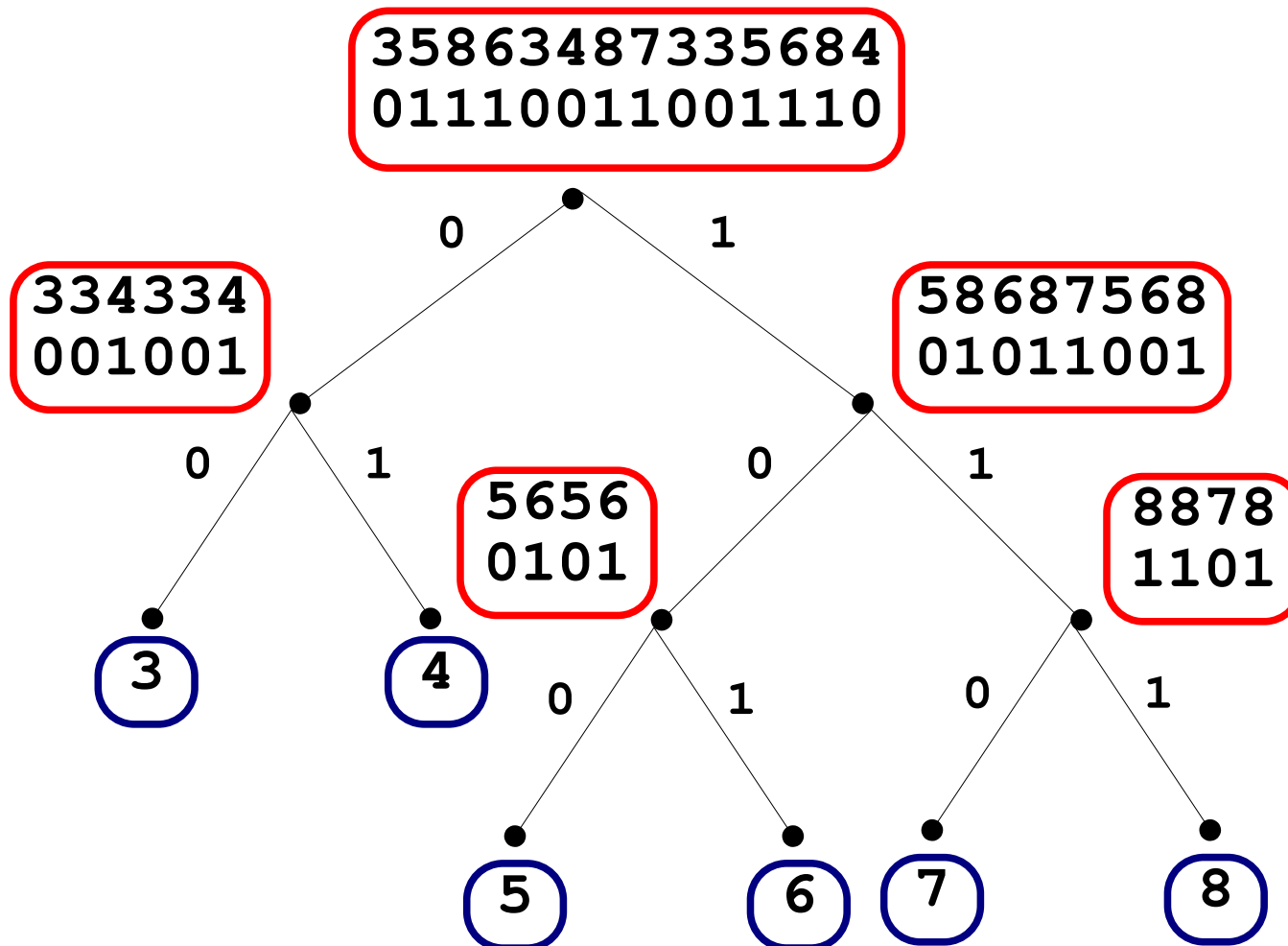
$\text{range_quantile}(S,i,j,k)$: return the k -th smallest value in $S[i\cdots j]$

$\text{range_next}(S,i,j,x)$: return the smallest value greater than x in $S[i\cdots j]$

$\text{range_count}(S,i,j)$: return # of distinct values in $S[i\cdots j]$

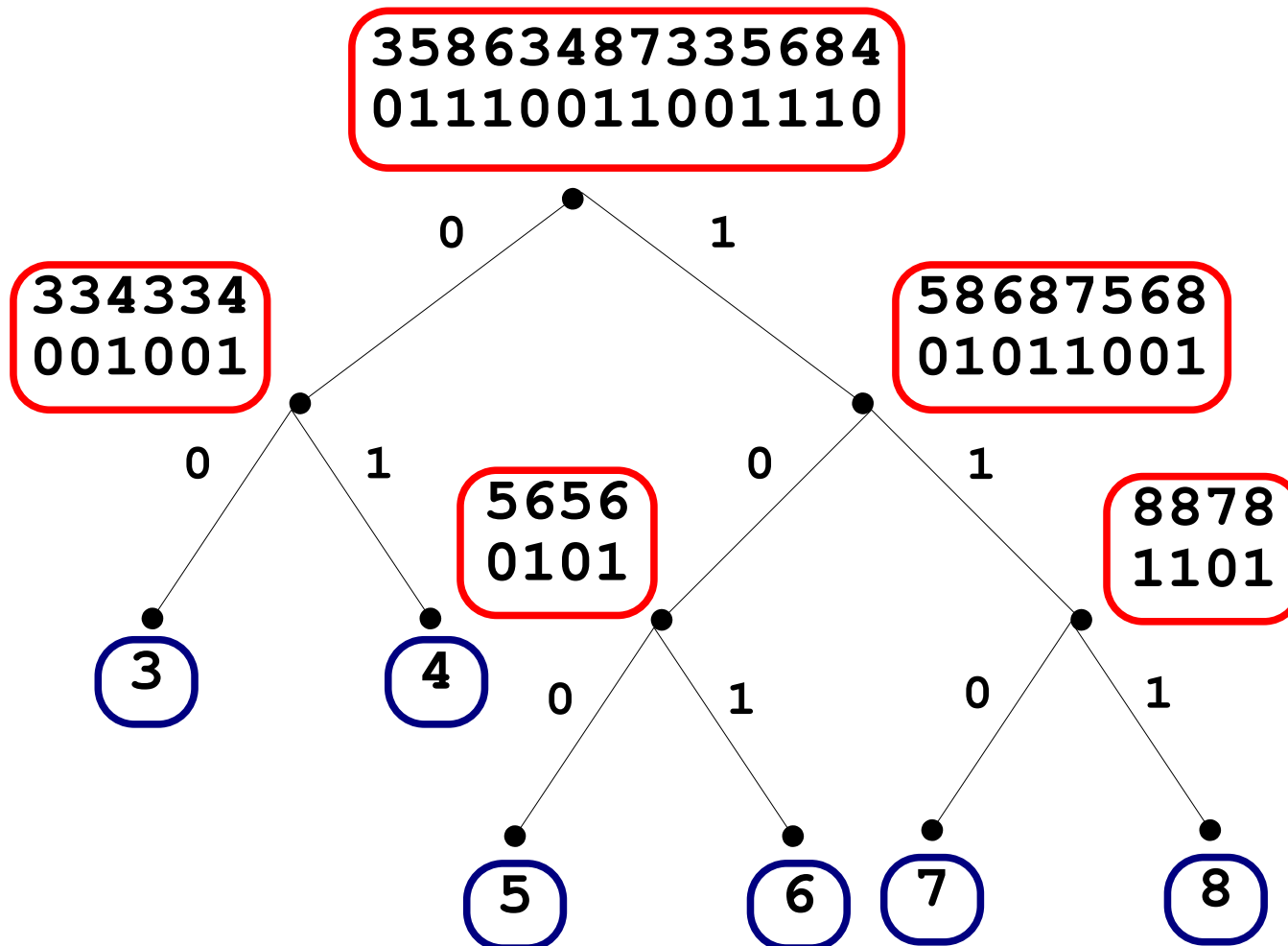
Example.

$S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]$



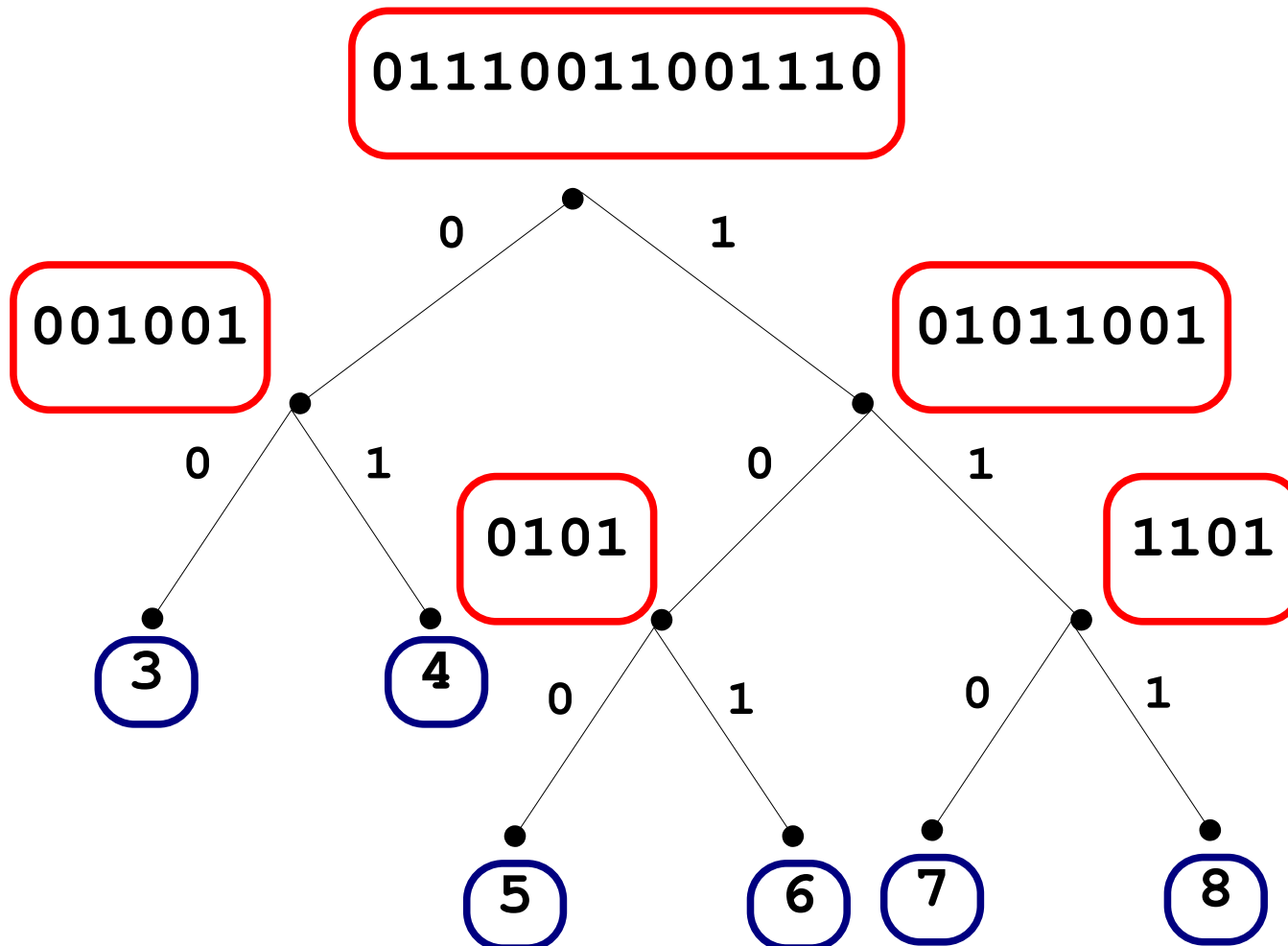
range_quantile(S,2,7,3): 3-rd smallest value in S[2...7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



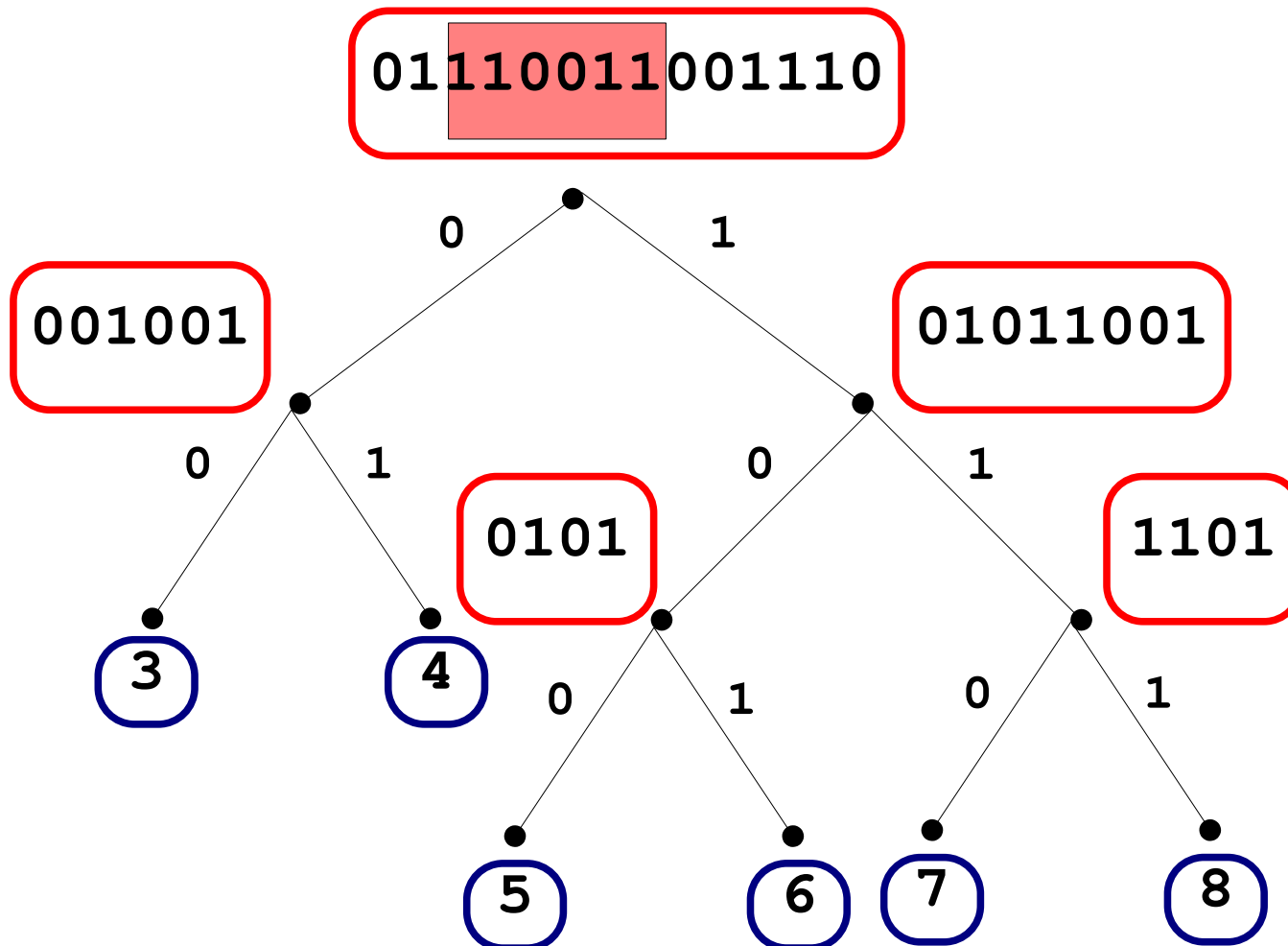
range_quantile(S,2,7,3): 3-rd smallest value in S[2...7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



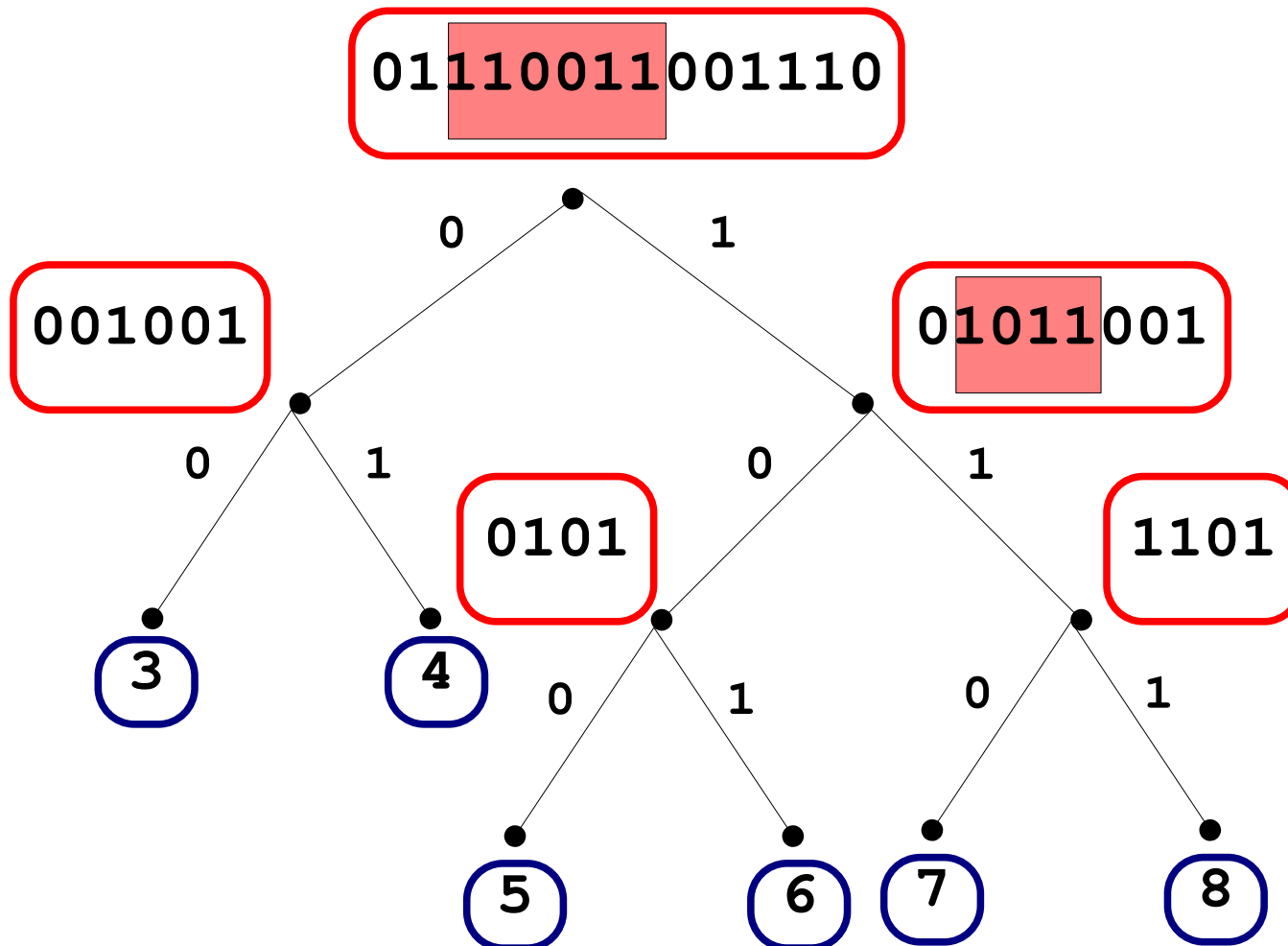
range_quantile(S,2,7,3): 3-rd smallest value in S[2..7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



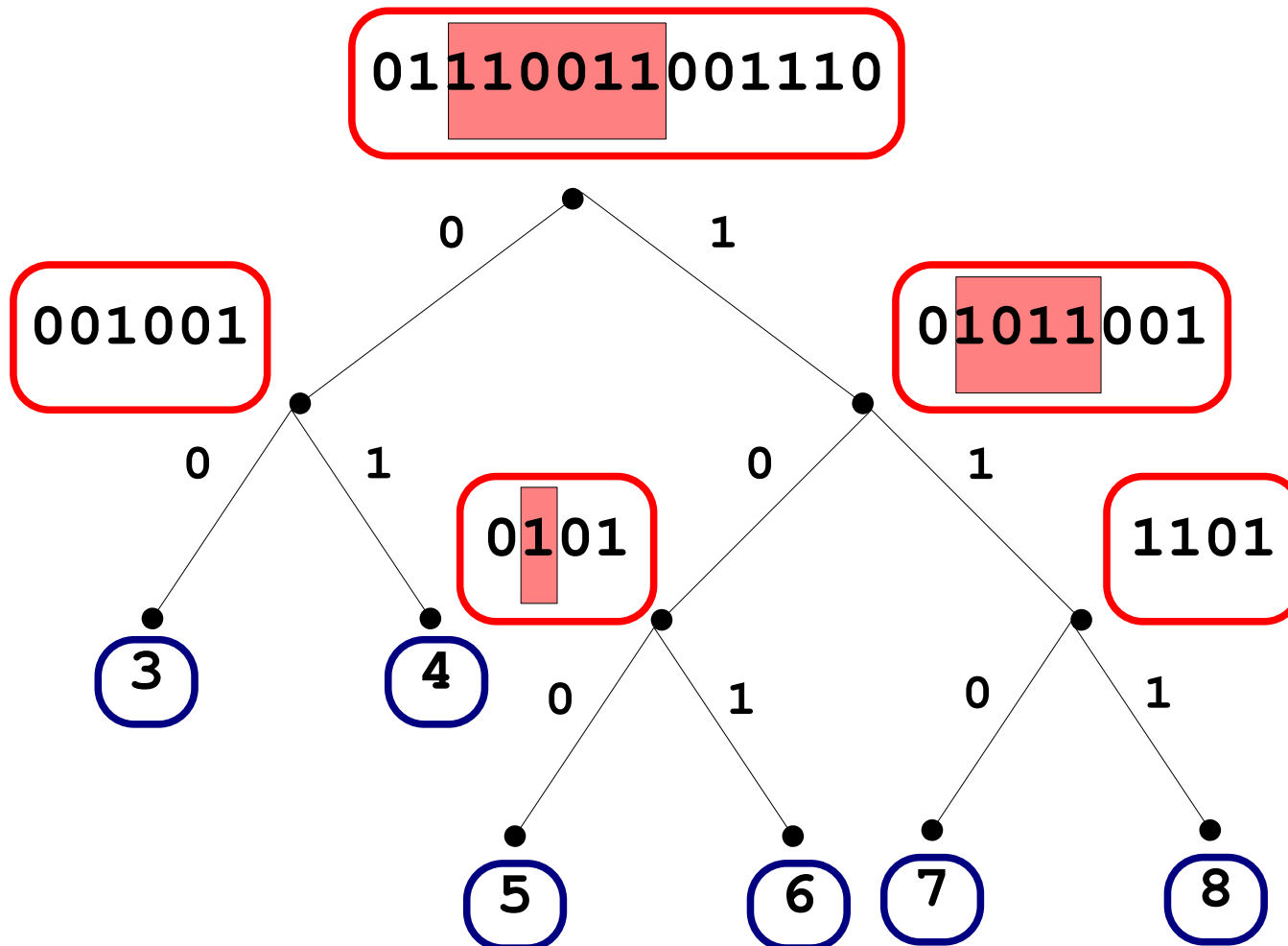
range_quantile(S,2,7,3): 3-rd smallest value in S[2...7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



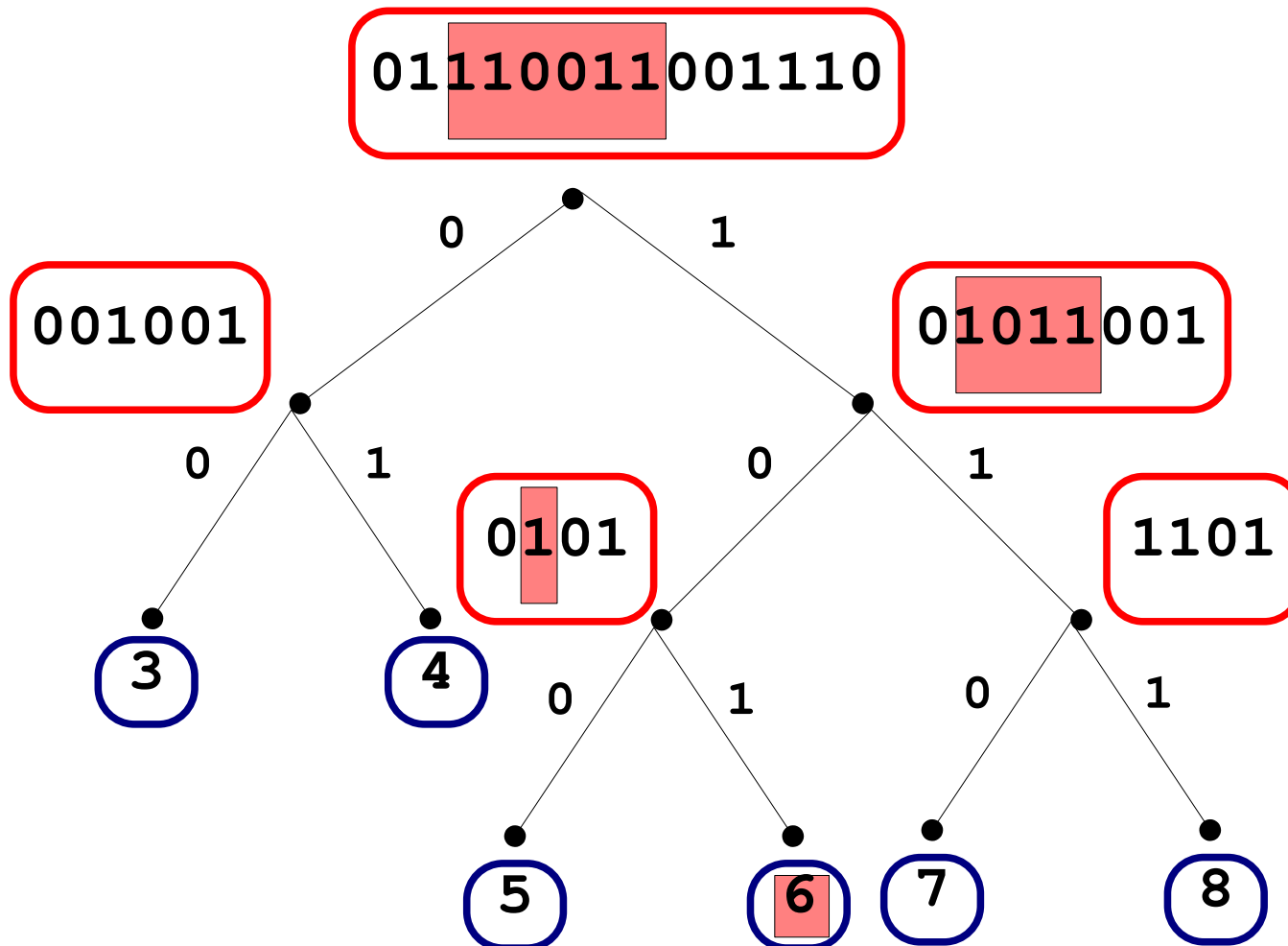
range_quantile(S,2,7,3): 3-rd smallest value in S[2..7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



range_quantile(S,2,7,3): 3-rd smallest value in S[2..7]

S = [3, 5, 8, 6, 3, 4, 8, 7, 3, 3, 5, 6, 8, 4]



Other applications

- Inverted indices (representation of posting lists)
- Computations Geometry (bidimensional range queries)
- Representation of Graphs, Permutations, ...
- Bioinformatics (maximal repeats)

Implementation

- Conceptually simple, but tricky in a few details
- Several implementations cited in the literature, some available on the web
- Ready to use, open source, modular library:

<https://github.com/simongog/sdsl-lite>

References

- R. Grossi, A. Gupta, J. Vitter [High-order entropy-compressed text indexes](#). [Proc. SODA 2003](#) (paper introducing WTs: somewhat hard to read)
- G. Navarro, V. Mäkinen. [Compressed Full-text indexes](#). [ACM Comp. Surv. 2007](#) (main reference for WT and compressed indices)
- G. Navarro. [Wavelet Trees for All](#), [Journal Discrete Algorithms](#), 2014 (recent review paper on WTs)

Thank you!